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ON 1-FACTORIZATIONS OF THE COMPLETE
GRAPH AND THE RELATIONSHIP TO
ROUND ROBIN SCHEDULES

ERIC NEIL GELLING

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GRAPH AND THE RELATIONSHIP TO
ROUND ROBIN SCHEDULES

by

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JULY 1973

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Supervisor: Dr. R. E. Odeh

ABSTRACT

The following new results concerning 1-factorizations of the complete graph are proved: (1) There are exactly 6 equivalence classes of 1-factorizations of the complete graph with 8 vertices. (2) There are exactly 396 equivalence classes of 1-factorizations of the complete graph with 10 vertices. Representatives of each of the equivalence classes are presented. The size of the automorphism group of each equivalence class of 1-factorizations of the complete graph with $2n$ vertices for $n \leq 5$ is also found.

Several theorems and results related to 1-factorizations of the complete graph are presented, and the relationship to round robin schedules is shown. An application problem demonstrates the importance of the choice of the equivalence class of round robin schedules in the solvability of the problem.

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CHAPTER 1

INTRODUCTION

In recent years the study of graph theory has been playing an increasingly important role in solving problems both mathematical and non-mathematical in nature. Up to about ten years ago, however, very little graph theory had been studied. Ore [17], in 1963, states in his bibliography that "the number of books on graph theory is very small" and then lists a mere nine of them, five of which are foreign works. Since he wrote those words a vast amount of material has been written on the subject. The computer has provided a powerful tool to graph theorists and much of the recent literature has been a result of the ability to use the computer to exhaustively generate and study characteristics of graphs. The potential for future study seems without limit.

It is the purpose of this thesis to study a particular area of graph theory: the 1-factorization of the complete graph with $2n$ vertices. The initial motivation for the study came from a specific problem of scheduling for a sporting event. It was desired to create a schedule with $2n$ teams such that every team plays every other team exactly once (known as a round robin). Furthermore, the games were to be played on n fields such that no team played more than twice on the same field nor twice in a row on the same

field. The problem is described in more detail in Chapter 4, along with a solution for 8 teams and 4 fields. In trying to solve the particular problem, the question arose whether all round robin schedules for $2n$ teams are the same except for re-labelling of teams and rounds. It was discovered that the answer is "yes" for $n \leq 3$ and "no" for $n \geq 4$.

In this study we have identified 6 equivalence classes for round robin schedules with 8 teams and 396 equivalence classes for round robin schedules with 10 teams. The problem is identical to the finding of the equivalence classes of 1-factorizations of the complete graph, and the relationship of the two problems is brought forth in Chapters 2 and 3.

In Chapter 2 we present a summary of graph theory beginning with the basic definitions and proceeding through some theorems which we use in the next chapter. Chapter 3 presents the main results of our study. Here we prove that there exists exactly 6 equivalence classes of 1-factorizations of the complete graph with 8 vertices; and exactly 396 equivalence classes of 1-factorizations of the complete graph with 10 vertices. In addition, we establish the size of the automorphism group for every equivalence class. To conclude the chapter we briefly present a method to produce a representative from each equivalence class of 1-factorizations of the complete graph with 12 vertices. As noted earlier, in Chapter 4 we solve the problem of scheduling for

a sporting event. In particular, we demonstrate that, when given certain constraints, the choice of the equivalence class of the round robin schedule can be vital as to whether the problem is solvable or not solvable.

In Appendix A we list representatives and characteristics of each of the equivalence classes of 1-factorizations of the complete graph with 10 vertices. Finally, in Appendix B we give a listing of the relevant computer programs.

CHAPTER 2

GRAPH THEORY

2-1. Introduction

It is the purpose of this section to introduce the terminology of graph theory that will be used later in the study. Most of the definitions are from Harary [8], but a few are from other works such as Wilson [24] and Behzad and Chartrand [2].

By a graph X we will mean a finite non-empty set V with p elements together with a finite set E of q unordered pairs of distinct elements of V . The elements of V are called vertices and the elements of E are called edges. We will frequently write $V(X)$ for the vertex set of the graph X and $E(X)$ for the edge set.

A labelling of the graph X with p vertices is a one-to-one mapping from $V(X)$ onto the set $\{0,1,2, \dots, p-1\}$. We will usually denote the vertex set $V(X)$ as $\{0,1,2, \dots, p-1\}$ and the edge set $E(X)$ as $\{[x,y]: x, y \in V(X), x \neq y\}$. The number of vertices in $V(X)$ (or the cardinality of $V(X)$) will be represented by $|V(X)|$. Similarly the cardinality of $E(X)$ will be $|E(X)|$.

If $e = [x,y] \in E(X)$ then e is incident with both x and y , and x and y are both incident with e . Two vertices (edges) joined by an edge (vertex) are said to be adjacent.

If X and Y are graphs such that $V(Y) \subseteq V(X)$ and $E(Y) \subseteq E(X)$ then $Y \subseteq X$ and Y is called a subgraph of X . If $V \subseteq V(X)$ then the subgraph $\langle V \rangle$ induced by V is the maximal subgraph of X whose vertex set is V ; i.e., it is the graph having vertex set V and whose edge set consists of those edges in X incident with two vertices of V . Similarly, if $E \subseteq E(X)$ then the subgraph $\langle E \rangle$ induced by E is the graph whose edge set is E and whose vertex set consists of those vertices of $V(X)$ incident with at least one edge of E . A subgraph of X is called a spanning subgraph if it contains all the vertices of X .

A walk of the graph X is an alternating sequence of vertices and edges $\{x_0, [x_0, x_1], x_1, [x_1, x_2], x_2, \dots, x_{m-1}, [x_{m-1}, x_m], x_m\}$ beginning and ending with vertices in which each edge is incident with the vertices immediately preceding and following it. In order to shorten the notation we will represent the walk by the sequence of vertices $\{x_0, x_1, \dots, x_m\}$ with the understanding that $[x_i, x_{i+1}]$ for $0 \leq i \leq m-1$ are the edges of the walk. A trail of the graph X is a walk such that the edges are distinct. A path of the graph X is a trail $\{x_0, x_1, \dots, x_m\}$ such that all vertices are distinct with the possible exception of x_0 and x_m . If $x_0 = x_m$ the path is closed and is called a circuit.

The order, k , of a circuit C is the number of edges (or vertices) in C . The circuit C , then, will be called a

k-circuit. It is clear that if the graph X has p vertices and $C \subseteq X$ then $3 \leq k \leq p$. If $C \subseteq X$ and the order of C is p then C spans the vertices of X and is called a Hamiltonian circuit. A graph with a Hamiltonian circuit is known as a Hamiltonian graph.

A graph is connected if every two vertices are joined by a path. A component of X is a maximal connected subgraph of X . Thus, a connected graph has only one component; a disconnected graph has more than one component.

Let X be a graph with p vertices with the condition that every pair of its p vertices is adjacent. Then X is called a complete p-graph and is denoted by K_p (after Kuratowski). We observe that the number of edges in K_p is given by:

$$|E(K_p)| = \frac{p(p-1)}{2} .$$

The complement \bar{X} of X is a graph which has $V(X)$ as the vertex set and in which two vertices are adjacent if and only if they are not adjacent in X , i.e.

$$V(\bar{X}) = V(X)$$

$$[x,y] \in E(\bar{X}) \text{ iff } x \neq y \text{ and } [x,y] \notin E(X) .$$

The degree of the vertex x is the number of edges incident with it. From this definition comes the first theorem of graph theory, known by Euler over two hundred years ago.

THEOREM 2.1. The sum of the degrees of the vertices of a graph X is twice the number of edges, i.e.,

$$\sum_{x \in V(X)} \deg(x) = 2|E(X)| .$$

This theorem is commonly called the "handshaking lemma" -- if several people shake hands, the total number of hands shaken must be even, precisely because there are always two hands to a shake. A natural corollary follows from this theorem.

COROLLARY 2.2. In any graph X, the number of vertices of odd degree is even.

If all vertices of X have the same degree, r, then X is called regular of degree r. A regular graph of degree 0 contains no edges; a regular graph of degree 1 has components consisting of exactly one edge. It is clear that the complete graph K_p is regular of degree p-1. It is also clear that the complement of a regular graph is regular.

2-2. 1-Factors, 1-Factorizations and Round Robin Schedules

In this section we present some definitions and theorems which will form a basis for the study of round robin schedules. A round robin schedule involving 2n teams has the property that every team plays every other team exactly

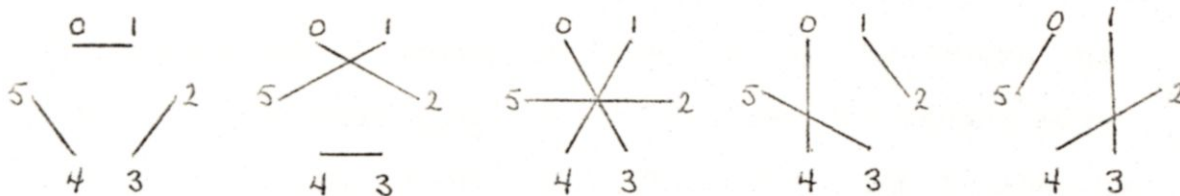
once. In a round robin schedule there are $2n-1$ rounds of n games each. We will define 1 -factors and a 1 -factorization of the complete graph K_{2n} and show the relationship to round robin schedules. In addition we shall present some theorems characteristic of 1 -factors of a graph.

A factor X_i of a graph X is a spanning subgraph of X which is not totally disconnected. The set of graphs $\{X_1, X_2, \dots, X_q\}$ is a decomposition of the graph X into the factors X_1, X_2, \dots, X_q if and only if the following are true:

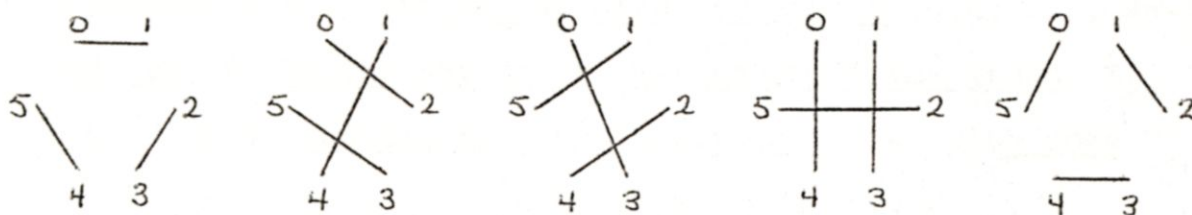
- i. $V(X_i) = V(X)$ for $1 \leq i \leq q$
- ii. $E(X_i) \cap E(X_j) = \phi$ for all $i \neq j$.
- iii. $E(X_1) \cup E(X_2) \cup \dots \cup E(X_q) = E(X)$.

An n -factor is a factor which is regular of degree n . If every X_i of the decomposition $\{X_1, X_2, \dots, X_q\}$ of the graph X is an n -factor then the decomposition is called an n -factorization, and X itself is n -factorable. Two examples of 1 -factorizations of K_6 are given in Figure 2.1.

A famous theorem of Tutte [21] characterizes graphs possessing a 1 -factor. Here, an odd component is a component of the graph X with an odd number of vertices.

FIGURE 2.1 TWO 1-FACTORIZATIONS OF K_6 

Example 1



Example 2

THEOREM 2.3. A graph X with p vertices has a 1-factor if and only if p is even and there is no subset $V \subseteq V(X)$ such that the number of odd components of $\langle V(X) - V \rangle$ exceeds $|V|$.

In this study we are interested in the 1-factorizations of the complete graph K_{2n} . Clearly (by observation or by Tutte's Theorem) K_{2n} contains a 1-factor with n edges.

THEOREM 2.4. The complete graph K_{2n} has $t = (2n-1)(2n-3) \dots (1)$ different 1-factors.

PROOF. Every 1-factor of K_{2n} covers the vertices and has n edges. By fixing one vertex of the $2n$ vertices, the first edge can be chosen in $(2n-1)$ ways. By fixing one vertex of the remaining $(2n-2)$ vertices the second edge can be chosen in $(2n-3)$ ways. The conclusion follows by induction on the n edges. ■

COROLLARY 2.5. Let the vertices of K_{2n} be $\{0,1, \dots, 2n-1\}$ and let W be the set of t 1-factors of the graph K_{2n} . Let $W_i \subseteq W$ contain all the 1-factors in W with edge $[0,i]$ for $i = 1, \dots, 2n-1$. Then W_1, \dots, W_{2n-1} represents a partition of W into $2n-1$ subsets with $s = (2n-3)(2n-5) \dots (1)$ members each.

PROOF. Follows from the proof of Theorem 2.4. ■

The fact that K_{2n} is 1-factorable is well known, and schemes for generating a 1-factorization (which we will show is equivalent to forming a round robin schedule) have been known since at least 1859 (Reisz [19]). Schemes are given in Ore [17], Moon [16], König [10], Freund [6], Lockwood [12] and [13], and Kraitichik [11]. The particular proof below is based on Harary [8].

THEOREM 2.6. The complete graph K_{2n} is 1-factorable.

PROOF. By definition, K_{2n} has $2n$ vertices and any 1-factor in K_{2n} has n edges. To be 1-factorable, the graph K_{2n} , which has $n(2n-1)$ edges, must be partitionable into $(2n-1)$ 1-factors. Let the vertices of K_{2n} be $\{0, 1, 2, \dots, 2n-1\}$. Let $B = \{1, 2, \dots, 2n-1\}$. Then for $\alpha \in B$ consider the edge-induced subgraph X_α of K_{2n} defined as follows:

$$E(X_\alpha) = \{[0, \alpha] \cup ([\alpha + \mu, \alpha - \mu] : \mu = 1, 2, \dots, n-1)\}$$

where each of the $\alpha + \mu, \alpha - \mu$ is an element of B modulo $(2n-1)$. A diagram of the subgraph is shown in Figure 2.2, where the vertices in B are on the circle, while the vertex 0 is outside the circle.

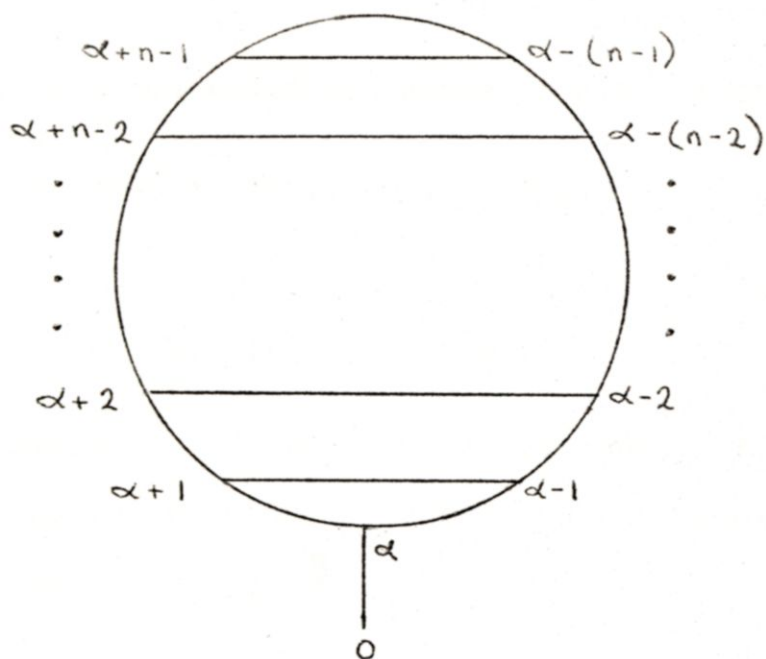
We claim that X_α is a 1-factor of K_{2n} . Clearly, since $0 \notin B$, the vertex 0 appears only once in $E(X_\alpha)$. Also, since $B - \{\alpha\}$ contains exactly $2(n-1)$ elements it is clear that the vertex α appears only once (i.e., in the edge $[0, \alpha]$) and that in fact every vertex appears exactly once. Therefore, X_α is a 1-factor for each $\alpha \in B$.

We claim that for $\beta \neq \alpha$ and $\beta, \alpha \in B$, X_β and X_α are edge disjoint. Clearly, $[0, \beta] \neq [0, \alpha]$. Assume

$$[\beta + \omega, \beta - \omega] = [\alpha + \mu, \alpha - \mu], \quad 1 \leq \mu, \omega \leq n-1$$

where each vertex is contained in B modulo $2n-1$.

FIGURE 2.2 A 1-FACTORIZATION OF K_{2n} BY THEOREM 2.6.



Without loss of generality, assume

$$\beta + \omega \equiv \alpha + \mu \pmod{2n-1}.$$

Then

$$\beta - \omega \equiv \alpha - \mu \pmod{2n-1}.$$

By adding we get

$$2\beta \equiv 2\alpha \pmod{2n-1}$$

and since $2n-1$ is odd we get

$$\beta \equiv \alpha \pmod{2n-1}.$$

But this is a contradiction since $\beta \neq \alpha$. Therefore

$$[\beta+\omega, \beta-\omega] \neq [\alpha+\mu, \alpha-\mu], \quad 1 \leq \mu, \omega \leq n-1$$

and

$$E(X_\alpha) \cap E(X_\beta) = \phi.$$

Since B contains $2n-1$ elements, $X_1 \cup X_2 \cup \dots \cup X_{2n-1}$ contains $n(2n-1)$ edges and represents a complete 1-factorization of K_{2n} . ■

A round robin schedule on $2n$ vertices (hereafter called simply a schedule) is a 1-factorization of the complete graph K_{2n} . The vertices of K_{2n} correspond to the teams, the edges correspond to the games, and the 1-factors correspond to the rounds. We shall use the terms "1-factorization" and "schedule" interchangeably in our study.

The following set of theorems and definitions characterize the nature of 1-factors in a graph.

THEOREM 2.7. Let X_i and X_j be two disjoint 1-factors of any graph X . Then every vertex of X has exactly one edge of X_i and exactly one edge of X_j incident with it.

PROOF. Since each 1-factor of X covers the vertices and since $E(X_i) \cap E(X_j) = \phi$, the theorem follows immediately. ■

In any graph X , an alternating trail is a trail such that the edges alternate between two disjoint 1-factors of X .

COROLLARY 2.8. Let X_i and X_j be two disjoint 1-factors of any graph X . Then, starting at any vertex of X , say x_0 , there exists an alternating trail which eventually returns to x_0 . Furthermore, the first vertex the alternating trail passes through twice is x_0 , and thus x_0 lies on a circuit.

PROOF. Follows immediately from Theorem 2.7. ■

In any graph X , an alternating circuit is a circuit such that the edges alternate between two disjoint 1-factors of X . If the alternating circuit covers the vertices of X then it is an alternating Hamiltonian circuit.

COROLLARY 2.9. Let X_i and X_j be two disjoint 1-factors of any graph X . Then every vertex of X is a member of an alternating circuit between X_i and X_j .

PROOF. It follows from Corollary 2.8 that every vertex lies on a circuit. But by Theorem 2.7, this circuit must be an alternating circuit. ■

It will be observed that, by the definition of an alternating circuit, the length k of an alternating circuit is even and furthermore $k \geq 4$.

COROLLARY 2.10. Let X_i and X_j be two disjoint 1-factors of any graph X . Then $X_i \cup X_j$ is the disjoint union of alternating circuits.

PROOF. Follows immediately from Corollary 2.9. ■

The next theorem and corollary are interesting results of Theorem 2.6. Let

$$F = \{Y_i ; 1 \leq i \leq 2n-1\}$$

be a 1-factorization of the complete graph K_{2n} . Then F is a completely Hamiltonian schedule (or an H-schedule) if, for every $Y_i, Y_j \in F$ and $i \neq j$, $Y_i \cup Y_j$ is an alternat-

ing Hamiltonian circuit. A completely non-Hamiltonian schedule (or an NH-schedule) is a schedule which has no Hamiltonian circuits between disjoint 1-factors.

THEOREM 2.11. Let K_{2n} be partitioned as in Theorem 2.6 and let X_α and $X_{\alpha+\beta}$ be two disjoint 1-factors defined by this partitioning. Then $X_\alpha \cup X_{\alpha+\beta}$ is an alternating Hamiltonian circuit if and only if there exists no integer c , $0 < c < 2n-1$, such that $c\beta = 0 \pmod{2n-1}$.

PROOF. X_α and $X_{\alpha+\beta}$ are defined in Theorem 2.6 as follows:

$$E(X_\alpha) = \{[0, \alpha] \cup ([\alpha+\mu, \alpha-\mu]: \mu = 1, 2, \dots, n-1)\}$$

$$E(X_{\alpha+\beta}) = \{[0, \alpha+\beta] \cup ([\alpha+\beta+\mu, \alpha+\beta-\mu]: \mu = 1, 2, \dots, n-1)\}$$

where each of the α , $\alpha+\beta$, $\alpha+\mu$, $\alpha-\mu$, $\alpha+\beta+\mu$, $\alpha+\beta-\mu$ is an element of $\{1, 2, \dots, 2n-1\}$ modulo $(2n-1)$. See Figure 2.3 for examples where $n = 5$ and $\beta = 3$ and 2.

By the definition of X_α and $X_{\alpha+\beta}$ and by Corollary 2.9 there exists an alternating circuit on $X_\alpha \cup X_{\alpha+\beta}$:

$$\{0, \alpha, \alpha+2\beta, \alpha-2\beta, \alpha+4\beta, \alpha-4\beta, \dots, \alpha+x\beta, \alpha-x\beta, 0\}.$$

Therefore for some even x , there exists an alternating circuit which includes the vertex o and which must contain $x+2$ vertices. However, since the edge $[0, \alpha+\beta] \in E(X_{\alpha+\beta})$

must be in the alternating circuit, we have

$$\alpha - x\beta = \alpha + \beta \pmod{2n-1}.$$

Let $c = x+1$. Then

$$\alpha - x\beta = \alpha - (c-1)\beta = \alpha + \beta \pmod{2n-1}$$

and hence

$$c\beta = 0 \pmod{2n-1}.$$

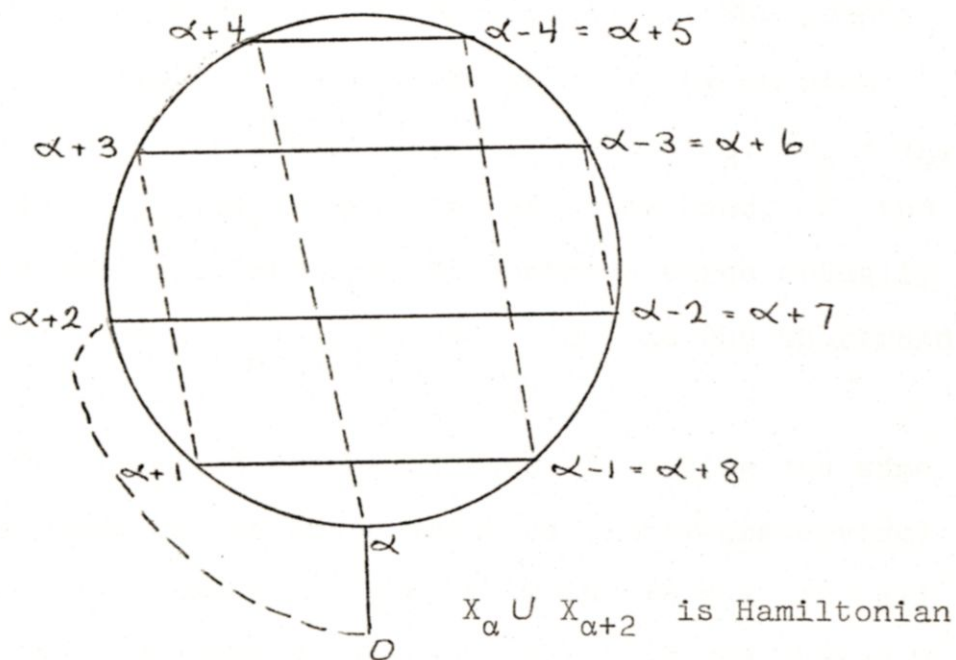
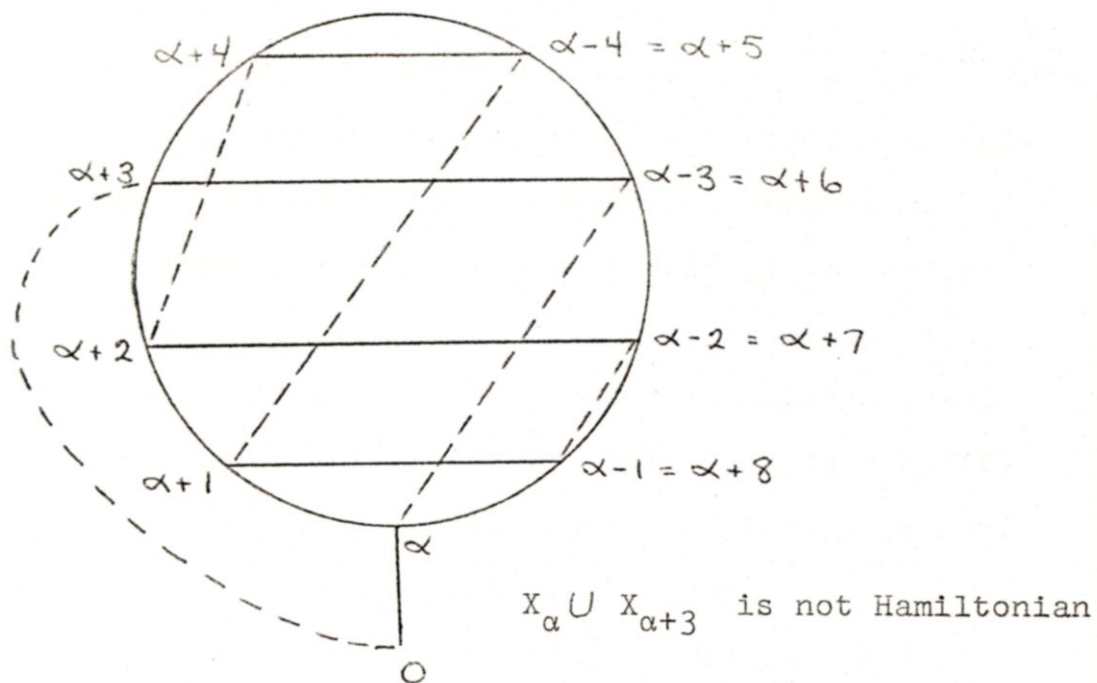
But since $X_\alpha \cup X_{\alpha+\beta}$ is an alternating Hamiltonian circuit if and only if $x+2 = 2n$, the conclusion follows. ■

COROLLARY 2.12. The schedule which is defined by the partitioning of K_{2n} in Theorem 2.6 is an H-schedule if and only if $2n-1$ is a prime number.

By Corollary 2.12 we observe that the schedules defined by Theorem 2.6 for $K_4, K_6, K_8, K_{12}, K_{14}, K_{18}, K_{20}, K_{24}$ must be H-schedules. Similarly, the schedules for $K_{10}, K_{16}, K_{22}, K_{26}$ must not be H-schedules.

An interesting unsolved question is whether or not for every $n \geq 2$ there exists an H-schedule and for every $n \geq 4$ whether or not there exists an NH-schedule.

FIGURE 2.3 ILLUSTRATION OF THEOREM 2.11 FOR 1-FACTORS

OF K_{10} .

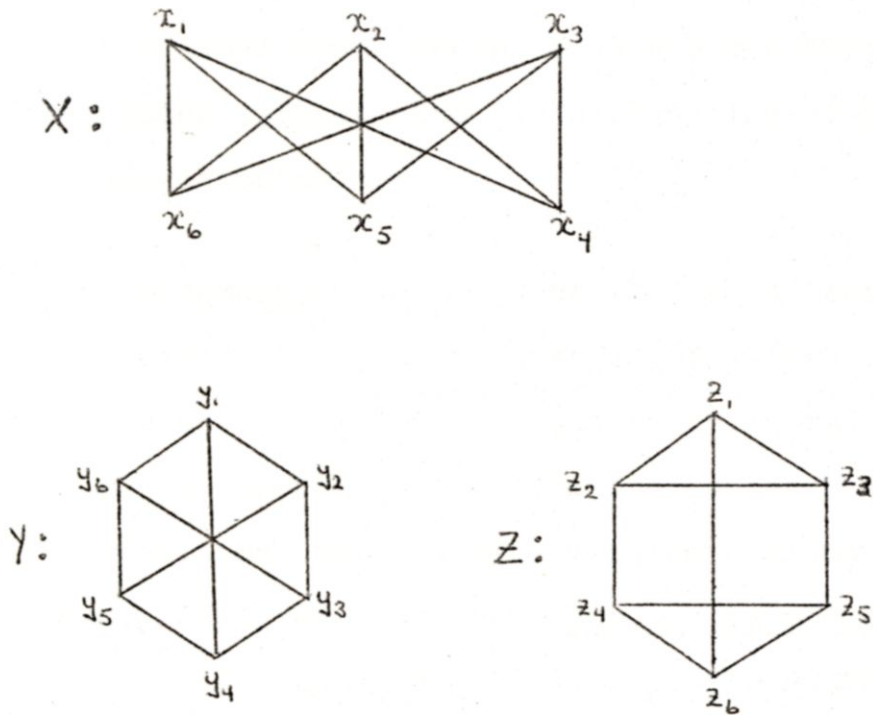
2-3. Isomorphisms, Automorphism Groups, and Equivalence Classes

Two graphs, X and Y , which may look quite different geometrically, may in fact be the same graph. For this purpose we shall introduce the concept of isomorphism.

Two graphs X and Y are isomorphic if there exists a one-to-one mapping ψ , called an isomorphism, from $V(X)$ onto $V(Y)$ such that ψ preserves adjacency and nonadjacency; i.e. $[x_1, x_2] \in E(X)$ iff $[\psi x_1, \psi x_2] \in E(Y)$. Clearly, isomorphism is an equivalence relation on graphs, and hence divides the collection of all graphs into equivalence classes. If there exists no mapping on $V(X)$ which will map X onto Y then X and Y are non-isomorphic and belong to different equivalence classes. The graphs X and Y of Figure 2.4 are isomorphic by the mapping $\psi : V(X) \rightarrow V(Y)$ defined by: $\psi x_1 = y_1$, $\psi x_2 = y_3$, $\psi x_3 = y_5$, $\psi x_4 = y_2$, $\psi x_5 = y_4$, $\psi x_6 = y_6$. On the other hand, X and Z are non-isomorphic because Z contains three mutually adjacent vertices but X does not. (Behzad and Chartrand [2], p. 3).

Two labelled graphs X and Y having the same labels are identical if there exists a (label-preserving) isomorphism ψ between X and Y with $\psi x = x$ for all $x \in V(X)$. By referring to Figure 2.4 it will be observed that X and Y are isomorphic but not identical.

FIGURE 2.4 AN ILLUSTRATION OF ISOMORPHISM AND NON-ISOMORPHISM.



The following two theorems are obvious results of the definitions.

THEOREM 2.13. If X_i and X_j are 1-factors of the graph X then X_i and X_j are isomorphic.

THEOREM 2.14. If C_i and C_j are two circuits of length k then C_i and C_j are isomorphic.

It follows from Theorem 2.14 and the definition of an alternating circuit that two alternating circuits of length k are isomorphic.

An automorphism of a graph X is an isomorphism of X with itself, i.e., a permutation of $V(X)$ which preserves adjacency. Clearly, an automorphism followed by another is an automorphism and therefore (under the operation of composition) the set of all automorphisms on X form a group, $G(X)$, known as the automorphism group of X . The order of $G(X)$, denoted $|G(X)|$, is the number of members in the group. Since the identity mapping on X is always in $G(X)$, $|G(X)| \geq 1$. The identity mapping is the trivial automorphism of X ; all other automorphisms, if any exist, are non-trivial. Using Tutte's definition [22, p. 52] a graph is symmetrical if it has a non-trivial automorphism.

THEOREM 2.15. For any graph X and its complement \bar{X} , $G(X)$ is isomorphic to $G(\bar{X})$.

PROOF. Every element ψ of $G(X)$ is a permutation on $V(X)$ which preserves adjacency in X . However, ψ preserves adjacency if and only if ψ preserves non-adjacency. Thus a permutation on $V(X)$ is an automorphism of X if and only if it is an automorphism on \bar{X} , implying that $G(X)$ is isomorphic to $G(\bar{X})$. ■

Several results follow from the above definitions.

Result 1: The order of K_p is $p!$ because a mapping from any vertex to any other vertex will preserve the structure of K_p .

Result 2: The order of a circuit of length k is $2k$. To show this let a circuit be represented by $C = \{x_0, x_1, \dots, x_k\}$ where $x_k = x_0$ and the subscripts are modulus k . Each subscript of the vertices can be increased by 1 without changing the circuit, thus transforming C into one of its rotations. Each subscript i can also be replaced by $k-i$, thus transforming C into its mirror image. Since there are k rotations and a mirror image of each, the order of C is $2k$.

Result 3: The order of a 1-factor with $2n$ vertices is $n!2^n$. To show this we observe that since there are n disjoint edges they can be permuted $n!$ ways. But since there are two vertices incident with each edge, there

are 2^n ways of permuting the vertices on the edges.

Result 4: The order of a 1-factor in a 1-factorization of K_{2n} is $(2n-1)n!2^n$. By result 3 above, the order of a 1-factor is $n!2^n$. But since a 1-factorization of K_{2n} has $(2n-1)$ 1-factors and (by Theorem 2.13) they are isomorphic, every 1-factor can be mapped onto every other 1-factor. Therefore, if X_i is a 1-factor in a 1-factorization of K_{2n} , there are $(2n-1)n!2^n$ mappings on the vertices of K_{2n} which will preserve X_i .

THEOREM 2.16. The number of ways in which a graph X with p vertices can be labelled is $\frac{p!}{|G(X)|}$.

PROOF. (Behzad and Chartrand [2], p. 167). Let $\{x_0, \dots, x_{p-1}\}$ be a set of p labels. Clearly there are $p!$ ways of labelling X without regard to the number of resulting labelled graphs which may be identical. Let X_1 and X_2 be two labelled graphs obtained from X . Then the relation " X_1 is identical to X_2 " is an equivalence relation on the set of labelled graphs obtained from X . For the labelled graph X_1 , an automorphism of X gives rise to a labelled graph identical to X_1 , and conversely. Hence, each equivalence class of a labelling X_1 of X contains $|G(X)|$ members. This implies that there are $p!/|G(X)|$ labellings in all. ■

Throughout this study we will let $L(X)$ represent the set of all labellings of X and $|L(X)|$ represent the number of elements in $L(X)$.

One of the more difficult problems of graph theory has been to find an easy test for isomorphism between graphs. Much research has been devoted by chemists, network analysts, linguists, logicians, etc., to solve this problem. Probably the most efficient procedure to date was developed by Corneil and Gotlieb in 1970 [5], but unfortunately it is based on a conjecture. It remains still unsolved to find an efficient deterministic algorithm (i.e. one which can be solved in time T where T is proportional to a constant power of the number of vertices of the graph) to determine whether two graphs are isomorphic.

CHAPTER 3

EQUIVALENCE CLASSES OF 1-FACTORIZATIONS
OF THE COMPLETE GRAPH K_{2n} 3-1. Introduction

In this chapter, vertices will be labelled from the set $\{1, 2, \dots, 2n\}$ in order to correspond to the usual labellings of teams in round robin schedules.

Throughout this chapter let the set W be the set of 1-factors of K_{2n} . For $i = 1, \dots, 2n-1$ we will let W_i represent the set of all 1-factors in W with the edge $[1, i+1]$. By Corollary 2.5, W will thus be partitioned into $2n-1$ sets, W_i for $i = 1, \dots, 2n-1$, each containing $s = (2n-3)(2n-5)\dots(1)$ 1-factors.

We observe that any 1-factorization of K_{2n} will have exactly one member from each W_i .

From Theorem 2.13 we know that all 1-factors of K_{2n} are isomorphic. Let $X_1, X_2 \in W_1$ such that $X_1 \neq X_2$. Let T_1 and T_2 be the complete sets of 1-factorizations of K_{2n} which contain X_1 and X_2 respectively. Then a mapping ψ which maps X_1 onto X_2 must also be a one-to-one mapping of T_1 onto T_2 . Therefore, in searching for equivalence classes of 1-factorizations of K_{2n} and in enumerating the members in each class, we can reduce our

search by fixing a particular 1-factor $X_1 \in W_1$ to be a member of every 1-factorization studied.

Without loss of generality, we will let X_1 be the 1-factor whose edge set is $\{[1,2], [3,4], \dots, [2n-1, 2n]\}$.

In this chapter we will study the 1-factorizations of K_{2n} for $n \leq 5$. We observe immediately that the 1-factorizations of K_2 and K_4 are trivial. The only 1-factors of K_4 are:

$$X_1: \{[1,2], [3,4]\}$$

$$X_2: \{[1,3], [2,4]\}$$

$$X_3: \{[1,4], [2,3]\}$$

Clearly, each W_i for $i = 1, 2, 3$ contains only one member and therefore $F = \{X_1, X_2, X_3\}$ must be the only 1-factorization of K_4 . By Theorem 2.16, the order $|G(F)|$ of the automorphism group of F is $4! = 24$.

By Corollary 2.5 the set of 1-factors, W , for K_6 contains 15 members and each W_i for $i = 1, \dots, 5$ contains 3 1-factors. They are as follows:

$$W_1 \left\{ \begin{array}{l} X_1: \{[1,2], [3,4], [5,6]\} \\ X_2: \{[1,2], [3,5], [4,6]\} \\ X_3: \{[1,2], [3,6], [4,5]\} \end{array} \right. \quad W_2 \left\{ \begin{array}{l} X_4: \{[1,3], [2,4], [5,6]\} \\ X_5: \{[1,3], [2,5], [4,6]\} \\ X_6: \{[1,3], [2,6], [4,5]\} \end{array} \right.$$

$$W_3 \begin{cases} X_7: \{[1,4], [2,3], [5,6]\} \\ X_8: \{[1,4], [2,5], [3,6]\} \\ X_9: \{[1,4], [2,6], [3,5]\} \end{cases} \quad W_4 \begin{cases} X_{10}: \{[1,5], [2,3], [4,6]\} \\ X_{11}: \{[1,5], [2,4], [3,6]\} \\ X_{12}: \{[1,5], [2,6], [3,4]\} \end{cases}$$

$$W_5 \begin{cases} X_{13}: \{[1,6], [2,3], [4,5]\} \\ X_{14}: \{[1,6], [2,4], [3,5]\} \\ X_{15}: \{[1,6], [2,5], [3,4]\} \end{cases}$$

It can easily be checked that there are only two 1-factorizations of K_6 which include the fixed 1-factor X_1 .

They are:

$$F_1: \{X_1, X_5, X_9, X_{11}, X_{13}\}$$

$$F_2: \{X_1, X_6, X_8, X_{10}, X_{14}\}$$

The mapping $\psi: F_2 \rightarrow F_1$ defined by:

$$1 \rightarrow 1, \quad 2 \rightarrow 2, \quad 3 \rightarrow 3, \quad 4 \rightarrow 4, \quad 5 \rightarrow 6, \quad 6 \rightarrow 5$$

shows that there is only one equivalence class, which we shall label $L(F_1)$. Clearly, since W_1 contains 3 members, $L(F_1)$ contains 6 members. By Theorem 2.16, the order of the automorphism group for $L(F_1)$ is $6!/6 = 120$.

3-2. Results for K_8 .

By Corollary 2.5 the set of 1-factors, W , for K_8 contains 105 members and each W_i for $i = 1, \dots, 7$ contains 15 1-factors. A computer program (see Computer Program 1 in Appendix B) was written to generate the set W and find all the 1-factorizations of K_8 . The program fixed the first 1-factor $X_1 \in W_1$ and chose 1-factors lexicographically out of each W_i for $i = 2, \dots, 7$, choosing only 1-factors which did not have a common edge with previously chosen 1-factors. The procedure followed is given by the flowchart and supporting explanation in Figure 3.1, and corresponds to the well-known method of efficient search known as backtrack programming. See Golomb and Baumert [7] or Walker [23] for a general description of backtrack programming.

FIGURE 3.1 FLOWCHART FOR GENERATING SCHEDULES FROM THE PARTITIONED SET OF 1-FACTORS OF K_8 .

Variables

- n = number of edges in a 1-factor.
 nm = $2n-1$ = number of 1-factors in a schedule.
 s = $(2n-3)(2n-5) \dots (1)$ = number of elements in each W_i .
 $W(i,k)$ = k^{th} element in the set W_i for $i = 1, \dots, nm$
and $k = 1, \dots, s$.

FIGURE 3.1 (cont'd)

$P(i)$ = current pointer to the 1-factor in each W_i for
 $i = 1, \dots, nm.$

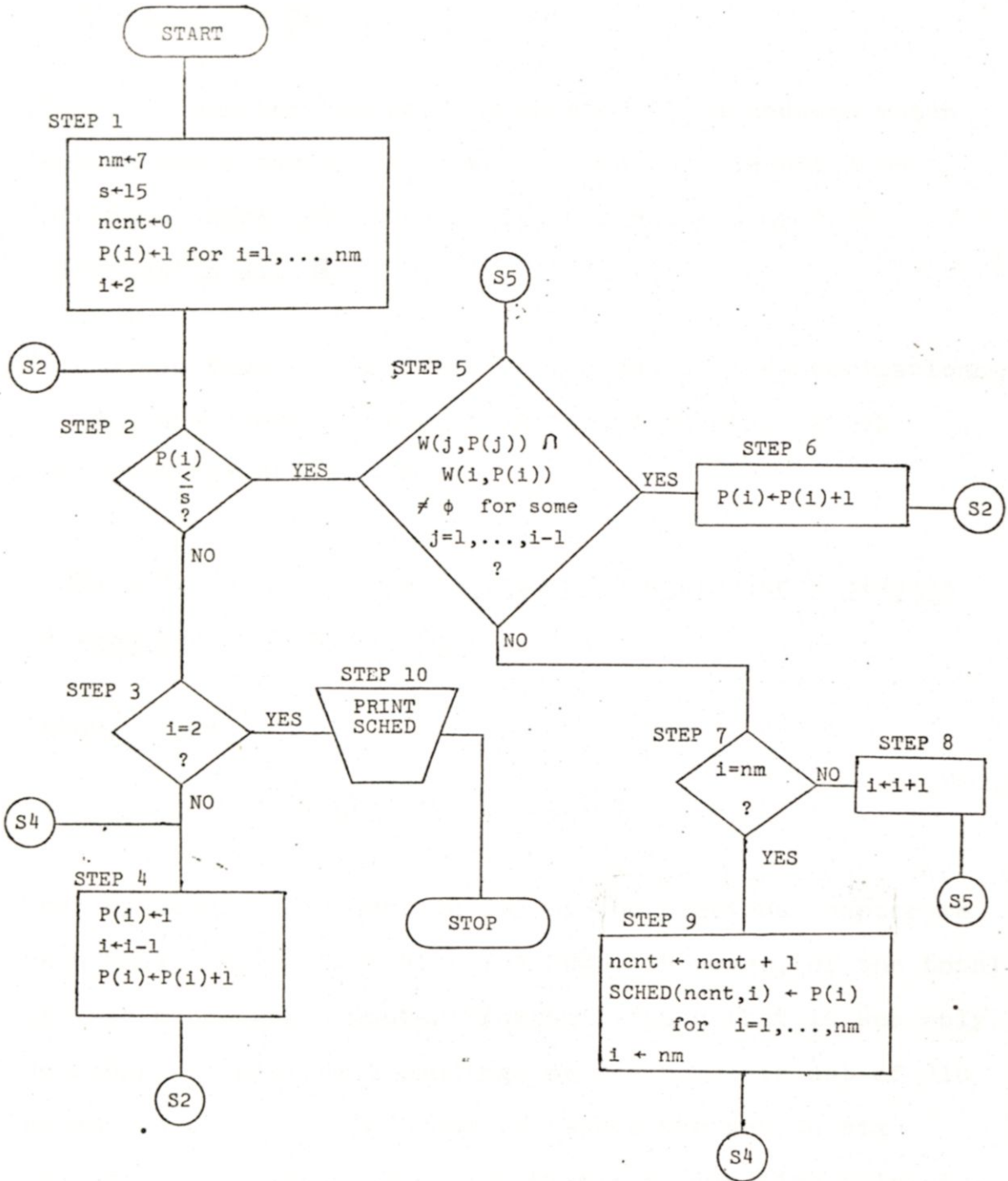
$ncnt$ = count of schedules.

$SCHED(ncnt, i)$ = output array for schedules.

Explanation of flowchart

- Step 1 - Initialize variables and arrays. (Note that i is initialized to 2 instead of to 1 because the program fixes the first factor in W_1).
- Step 2 - Are there any 1-factors left in W_i which are candidates? If yes, go to Step 5.
- Step 3 - Are we finished, i.e. are there any candidates left in W_2 ? If yes, go to Step 10.
- Step 4 - Initialize pointer (P_i) in W_i to 1 and backtrack to previous W_i . Pick next candidate in new W_i . Go to Step 2.
- Step 5 - Does the current candidate from W_i have edges in common with candidates from W_j , $j = 1, \dots, i-1$? If no, go to Step 7.
- Step 6 - Choose next candidate in W_i and go to Step 2.
- Step 7 - Has a schedule been found? If yes, go to Step 9.
- Step 8 - Move forward to next W_i and go to Step 5.
- Step 9 - Move schedule to output array. Set i to last W_i in preparation for backtrack and go to Step 4.
- Step 10 - Print all schedules and stop.

FIGURE 3.1 (cont'd)
FLOWCHART



THEOREM 3.1. There are 6240 1-factorizations of K_8 labelled on the set $\{1,2, \dots, 8\}$.

PROOF. Computer Program 1 generated 416 schedules which include the 1-factor $X_1 \in W_1$. Since W_1 contains 15 1-factors there are 15 sets of 416 schedules or 6240 schedules in all. ■

To find the equivalence classes of 1-factorizations of K_8 the computer program performed mappings which preserved the fixed 1-factor X_1 .

LEMMA 3.2. There are a maximum of 6 equivalence classes of 1-factorizations of K_8 .

PROOF. Let

$$G_1 = \{\psi \in G(K_8) : \psi X_1 = X_1\} .$$

Then by Result 4 in Section 2-3 of the previous chapter we know that $|G_1| = 2^4 \cdot 4! \cdot 7 = 2688$. However, of the total of 2688 mappings, Computer Program 1 found that it was only necessary to perform 9 mappings on the complete set of 416 schedules with X_1 in order to reduce the set to six equivalence classes. The nine mappings (ignoring trivial mappings of a vertex onto itself) are:

- $\psi_1: 1 \rightarrow 2, 2 \rightarrow 1$
 $\psi_2: 3 \rightarrow 4, 4 \rightarrow 3$
 $\psi_3: 5 \rightarrow 6, 6 \rightarrow 5$
 $\psi_4: 7 \rightarrow 8, 8 \rightarrow 7$
 $\psi_5: 5 \rightarrow 7, 7 \rightarrow 5, 6 \rightarrow 8, 8 \rightarrow 6$
 $\psi_6: 3 \rightarrow 5, 5 \rightarrow 3, 4 \rightarrow 6, 6 \rightarrow 4$
 $\psi_7: 1 \rightarrow 3, 3 \rightarrow 1, 2 \rightarrow 4, 4 \rightarrow 2$
 $\psi_8: 2 \rightarrow 3, 3 \rightarrow 2, 5 \rightarrow 8, 8 \rightarrow 5$
 $\psi_9: 2 \rightarrow 6, 3 \rightarrow 2, 6 \rightarrow 7, 7 \rightarrow 3$ ■

Figure 3.2 gives a listing of representatives of the six equivalence classes found by Computer Program 1.

FIGURE 3.2 REPRESENTATIVES FROM EACH OF THE EQUIVALENCE CLASSES OF 1-FACTORIZATIONS OF K_8 .

- $F_1: \{X_1, X_2, X_4, X_7, X_{11}, X_{16}, X_{22}\}$
 $F_2: \{X_1, X_2, X_4, X_7, X_{11}, X_{17}, X_{21}\}$
 $F_3: \{X_1, X_2, X_4, X_{10}, X_{13}, X_{15}, X_{20}\}$
 $F_4: \{X_1, X_2, X_4, X_8, X_{12}, X_{17}, X_{19}\}$
 $F_5: \{X_1, X_2, X_5, X_9, X_{13}, X_{14}, X_{20}\}$
 $F_6: \{X_1, X_3, X_6, X_9, X_{13}, X_{14}, X_{18}\}$

FIGURE 3.2 (cont'd)

The 1-factors X_1, X_2, \dots, X_{22} are as follows:

$$X_1: \{[1,2], [3,4], [5,6], [7,8]\}$$

$$X_2: \{[1,3], [2,4], [5,7], [6,8]\}$$

$$X_3: \{[1,3], [2,5], [4,7], [6,8]\}$$

$$X_4: \{[1,4], [2,3], [5,8], [6,7]\}$$

$$X_5: \{[1,4], [2,5], [3,8], [6,7]\}$$

$$X_6: \{[1,4], [2,6], [3,8], [5,7]\}$$

$$X_7: \{[1,5], [2,6], [3,7], [4,8]\}$$

$$X_8: \{[1,5], [2,6], [3,8], [4,7]\}$$

$$X_9: \{[1,5], [2,7], [3,6], [4,8]\}$$

$$X_{10}: \{[1,5], [2,7], [3,8], [4,6]\}$$

$$X_{11}: \{[1,6], [2,5], [3,8], [4,7]\}$$

$$X_{12}: \{[1,6], [2,7], [3,5], [4,8]\}$$

$$X_{13}: \{[1,6], [2,8], [3,7], [4,5]\}$$

$$X_{14}: \{[1,7], [2,3], [4,6], [5,8]\}$$

$$X_{15}: \{[1,7], [2,5], [3,6], [4,8]\}$$

$$X_{16}: \{[1,7], [2,8], [3,5], [4,6]\}$$

$$X_{17}: \{[1,7], [2,8], [3,6], [4,5]\}$$

$$X_{18}: \{[1,8], [2,4], [3,5], [6,7]\}$$

FIGURE 3.2 (cont'd)

X_{19} : {[1,8], [2,5], [3,7], [4,6]}

X_{20} : {[1,8], [2,6], [3,5], [4,7]}

X_{21} : {[1,8], [2,7], [3,5], [4,6]}

X_{22} : {[1,8], [2,7], [3,6], [4,5]}

THEOREM 3.3. There are exactly 6 equivalence classes of 1-factorizations of K_8 .

PROOF. Lemma 3.2 proves that there are no more than 6 equivalence classes. However, a count of the number of alternating Hamiltonian circuits in each of the representatives F_1, \dots, F_6 in Figure 3.2 shows that there are 0, 8, 12, 14, 18 and 21 alternating Hamiltonian circuits in the F_1, \dots, F_6 respectively. Clearly, therefore, they have different structure and they are non-equivalent. ■

A natural corollary follows immediately from this theorem.

COROLLARY 3.4. The number of alternating Hamiltonian circuits in any 1-factorization of K_8 will characterize the 1-factorization with respect to its equivalence class.

THEOREM 3.5. The number of members in each equivalence class of 1-factorizations of K_8 and the order of the automorphism groups are as follows:

| <u>Class</u> | <u>$L(F_i)$</u> | <u>$G(F_i)$</u> |
|--------------|------------------------------|------------------------------|
| $L(F_1)$ | 30 | 1344 |
| $L(F_2)$ | 630 | 64 |
| $L(F_3)$ | 420 | 96 |
| $L(F_4)$ | 2520 | 16 |
| $L(F_5)$ | 1680 | 24 |
| $L(F_6)$ | 960 | 42 |

where the F_i , $i = 1, \dots, 6$, correspond to the 1-factorizations listed in Figure 3.2.

PROOF. The count of the number of members in each $L(F_i)$ with the fixed 1-factor $X_1 \in W_1$ was performed by the computer program. Since there are 15 1-factors in W_1 each of the computer generated counts was multiplied by 15 in order to give the total number, $|L(F_i)|$, of members in each $L(F_i)$. By Theorem 2.16, $|G(F_i)| = 8!/|L(F_i)|$. ■

3-3. Results for K_{10}

By Corollary 2.5, the set of 1-factors, W , for K_{10} contains 945 members and each W_i for $i = 1, \dots, 9$ contains 105 1-factors. It was found that Computer Program 1 would use up far too much computer time to find all the 1-factorizations of K_{10} . It became necessary to reduce the search for schedules by reducing both the number of 1-factors to be considered and to put restrictions on the searches in the backtracking technique. The following set of lemmas enabled us to solve the problem.

LEMMA 3.6. For K_{10} , let W'_2 be the set of all 1-factors $X_1 \in W_2$ such that $E(X_1) \cap E(X_1) = \phi$. Then $X_1 \cup X_1$ is isomorphic to either $X_1 \cup X_2$ or $X_1 \cup X_3$ where X_2 and X_3 are given in Figure 3.3. There are 20 members in W'_2 which, in union with X_1 , are isomorphic to $X_1 \cup X_2$. There are 48 members in W'_2 which, in union with X_1 , are isomorphic to $X_1 \cup X_3$.

PROOF. Of the 105 1-factors in W_2 , only 68 do not have an edge in common with X_1 and therefore W'_2 has 68 members. By Corollary 2.10, $X_1 \cup X_1$ for $X_1 \in W'_2$ is the disjoint union of alternating circuits. Clearly $X_1 \cup X_1$ can only be a Hamiltonian circuit or the disjoint union of a

4-circuit and a 6-circuit. But by Theorem 2.14, all k -circuits are isomorphic. Therefore for every $X_i \in W_2'$, $X_1 \cup X_i$ must be isomorphic to either $X_1 \cup X_2$ (which is a 4-circuit and a 6-circuit) or $X_1 \cup X_3$ (which is a Hamiltonian circuit). By testing all 68 members in W_2' , it was found that 48 of them form a Hamiltonian circuit with X_1 and the remaining 20 form a 4-circuit and a 6-circuit. ■

FIGURE 3.3 REPRESENTATIVE 1-FACTORS OF K_{10} WITH MULTIPLICATION FACTORS FOR FINDING SIZES OF EQUIVALENCE CLASSES OF 1-FACTORIZATIONS.

Key

- M1 # 1-factors in W_1 isomorphic to X_1 .
- M2 # 1-factors in W_2 which, in union with X_1 , are isomorphic to $X_1 \cup X_i$ for $i = 2, 3$.
- M3 # 1-factors in W_3 which, in union with $X_1 \cup X_2$, are isomorphic to $X_1 \cup X_2 \cup X_i$ for $i = 5, 7, 8, \dots, 12$ such that X_1 and $X_1 \cup X_2$ are preserved.
- M4 # 1-factors in W_3 which, in union with $X_1 \cup X_3$, are isomorphic to $X_1 \cup X_3 \cup X_i$ for $i = 4, \dots, 7, 13, \dots, 36$ such that X_1 and $X_1 \cup X_3$ are preserved.

FIGURE 3.3 (cont'd)

| <u>1-Factor</u> | <u>Edges [x,y]</u> | <u>M1</u> | <u>M2</u> | <u>M3</u> | <u>M4</u> |
|-----------------|--------------------------------------|-----------|-----------|-----------|-----------|
| X ₁ | {[1,2], [3,4], [5,6], [7,8], [9,10]} | 105 | | | |
| X ₂ | {[1,3], [2,4], [5,7], [6,9], [8,10]} | | 20 | | |
| X ₃ | {[1,3], [2,5], [4,7], [6,9], [8,10]} | | 48 | | |
| X ₄ | {[1,4], [2,3], [5,7], [6,10], [8,9]} | | | | 2 |
| X ₅ | {[1,4], [2,3], [5,8], [6,10], [7,9]} | | | 3 | 3 |
| X ₆ | {[1,4], [2,3], [5,9], [6,8], [7,10]} | | | | 2 |
| X ₇ | {[1,4], [2,3], [5,10], [6,8], [7,9]} | | | 1 | 1 |
| X ₈ | {[1,4], [2,5], [3,6], [7,10], [8,9]} | | | 6 | |
| X ₉ | {[1,4], [2,5], [3,7], [6,10], [8,9]} | | | 6 | |
| X ₁₀ | {[1,4], [2,5], [3,8], [6,10], [7,9]} | | | 12 | |
| X ₁₁ | {[1,4], [2,5], [3,10], [6,7], [8,9]} | | | 6 | |
| X ₁₂ | {[1,4], [2,5], [3,10], [6,8], [7,9]} | | | 6 | |
| X ₁₃ | {[1,4], [2,6], [3,7], [5,10], [8,9]} | | | | 2 |
| X ₁₄ | {[1,4], [2,6], [3,8], [5,10], [7,9]} | | | | 1 |
| X ₁₅ | {[1,4], [2,6], [3,9], [5,8], [7,10]} | | | | 3 |
| X ₁₆ | {[1,4], [2,6], [3,10], [5,8], [7,9]} | | | | 1 |
| X ₁₇ | {[1,4], [2,7], [3,5], [6,10], [8,9]} | | | | 3 |
| X ₁₈ | {[1,4], [2,7], [3,6], [5,10], [8,9]} | | | | 2 |

FIGURE 3.3 (cont'd)

| <u>1-Factor</u> | <u>Edges [x,y]</u> | <u>M1</u> | <u>M2</u> | <u>M3</u> | <u>M4</u> |
|-----------------|--------------------------------------|-----------|-----------|-----------|-----------|
| X ₁₉ | {[1,4], [2,7], [3,9], [5,8], [6,10]} | | | | 1 |
| X ₂₀ | {[1,4], [2,7], [3,9], [5,10], [6,8]} | | | | 1 |
| X ₂₁ | {[1,4], [2,7], [3,10], [5,9], [6,8]} | | | | 2 |
| X ₂₂ | {[1,4], [2,8], [3,5], [6,10], [7,9]} | | | | 2 |
| X ₂₃ | {[1,4], [2,8], [3,6], [5,10], [7,9]} | | | | 1 |
| X ₂₄ | {[1,4], [2,8], [3,9], [5,7], [6,10]} | | | | 1 |
| X ₂₅ | {[1,4], [2,8], [3,9], [5,10], [6,7]} | | | | 1 |
| X ₂₆ | {[1,4], [2,8], [3,10], [5,9], [6,7]} | | | | 2 |
| X ₂₇ | {[1,4], [2,9], [3,5], [6,8], [7,10]} | | | | 1 |
| X ₂₈ | {[1,4], [2,9], [3,6], [5,8], [7,10]} | | | | 1 |
| X ₂₉ | {[1,4], [2,9], [3,7], [5,10], [6,8]} | | | | 1 |
| X ₃₀ | {[1,4], [2,9], [3,8], [5,10], [6,7]} | | | | 1 |
| X ₃₁ | {[1,4], [2,9], [3,10], [5,7], [6,8]} | | | | 1 |
| X ₃₂ | {[1,4], [2,9], [3,10], [5,8], [6,7]} | | | | 1 |
| X ₃₃ | {[1,4], [2,10], [3,5], [6,8], [7,9]} | | | | 1 |
| X ₃₄ | {[1,4], [2,10], [3,6], [5,8], [7,9]} | | | | 1 |
| X ₃₅ | {[1,4], [2,10], [3,9], [5,7], [6,8]} | | | | 1 |
| X ₃₆ | {[1,4], [2,10], [3,9], [5,8], [6,7]} | | | | 1 |

LEMMA 3.7. For K_{10} , let W'_3 be the set of 1-factors X_j in W_3 such that $E(X_1 \cup X_2) \cap E(X_j) = \phi$, where X_1 and X_2 are given in Figure 3.3. Then W'_3 has 40 members. By performing mappings which are automorphisms on both X_1 and $X_1 \cup X_2$, the set of $X_1 \cup X_2 \cup X_j$ for $X_j \in W'_3$ can be reduced to the 7 equivalence classes given in Figure 3.3 with entries in column M3. The numbers in column M3 give the number of members of W'_3 in each equivalence class.

PROOF. By exhaustion using Computer Program 2. ■

LEMMA 3.8. For K_{10} , let W^2_3 be the set of 1-factors X_j in W_3 such that $E(X_1 \cup X_3) \cap E(X_j) = \phi$, where X_1 and X_3 are given in Figure 3.3. Then W^2_3 has 41 members. By performing mappings which are automorphisms on X_1 and $X_1 \cup X_3$ the set of $X_1 \cup X_3 \cup X_j$ for $X_j \in W^2_3$ can be reduced to the 28 equivalence classes given in Figure 3.3 with entries in column M4. The numbers in column M4 give the number of members of W^2_3 in each equivalence class.

PROOF. By exhaustion using Computer Program 2. ■

We shall now introduce some notation which we shall use throughout the rest of the chapter.

Let

$$F = \{Y_1, \dots, Y_{2n-1}\}$$

represent a 1-factorization of K_{2n} such that each $Y_i \in W_i$. Then an alternating k -circuit in F is a circuit of length k such that the edges alternate between two 1-factors of F .

Let

$$A_k(F) = \{C : C \text{ is an alternating } k\text{-circuit for } F\}.$$

For $Y \in F$, let

$$A_k(F; Y) = \{C : C \text{ is an alternating } k\text{-circuit such that } E(C) \cap Y \neq \emptyset\}.$$

Let

$$P_k(F) = \{|A_k(F; Y)| : Y \in F\}.$$

For the vertex $x \in V(K_{2n})$ let

$$A_k(F; x) = \{C : C \text{ is an alternating } k\text{-circuit such that } x \in V(C)\}.$$

Let

$$Q_k(F) = \{|A_k(F; x)| : x \in V(K_{2n})\}.$$

The following lemma is useful for testing for isomorphism between two 1-factorizations of K_{2n} .

LEMMA 3.9. Let F and F' be two equivalent 1-factorizations of K_{2n} and let $\psi \in G(K_{2n})$ such that $\psi F = F'$.

Then:

- i. For $Y \in F$ and $Y' \in F'$, if $\psi Y = Y'$ then
 $|A_k(F; Y)| = |A_k(F'; Y')|$ for all k .
- ii. For $x, x' \in V(K_{2n})$, if $\psi x = x'$ then
 $|A_k(F; x)| = |A_k(F'; x')|$ for all k .
- iii. For $Y_i, Y_j \in F$ and $Y'_i, Y'_j \in F'$, if $\psi(Y_i \cup Y_j) =$
 $Y'_i \cup Y'_j$ then $\{|A_k(F; Y_i)|, |A_k(F; Y_j)|\} =$
 $\{|A_k(F', Y'_i)|, |A_k(F', Y'_j)|\}$ for all k .

PROOF. Follows immediately from the fact that isomorphism preserves adjacency and non-adjacency. ■

The purpose of the next lemma is to enable us to determine the number of members in an equivalence class $L(F)$ without generating all the members of $L(F)$.

LEMMA 3.10. Let F be defined as before and let

$$M_{k,i}(F) = \{F' : F' \in L(F) \text{ and } |A_k(F; Y_1)| = |A_k(F'; Y'_i)|\} \text{ for } i = 1, \dots, 2n-1.$$

Then

$$|M_{k,1}(F)| = \dots = |M_{k,2n-1}(F)| = M_k.$$

Also,

$$|L(F)| = \frac{M_k \cdot (2n-1)}{c}$$

where c ($c > 0$) is the multiplicity of $|A_k(F, Y_1)|$ in $P_k(F)$.

PROOF. By Lemma 3.9 every $F' \in L(F)$ is a member of some $M_{k,i}(F)$. Since by Theorem 2.13 all 1-factors of K_{2n} are isomorphic, it is obvious that

$$|M_{k,1}(F)| = \dots = |M_{k,2n-1}(F)|.$$

Let $M_k = |M_{k,1}(F)|$. If $|A_k(F, Y_1)|$ has multiplicity c in $P_k(F)$ then any $F' \in L(F)$ must be a member of c sets of $M_{k,i}(F)$ for $i = 1, \dots, 2n-1$.

Therefore

$$|L(F)| = \frac{M_k \cdot (2n-1)}{c} . \quad \blacksquare$$

LEMMA 3.11. For any 1-factorization F of K_{2n} let $A_k(F; Y_i)$ and $P_k(F)$ be defined as above. Then at least one $|A_k(F; Y_i)| \in P_k(F)$ is even.

PROOF. Let $|A_k(F)|$ represent the total number of alternating k-circuits in F. But every alternating circuit involves two 1-factors. Therefore, by the definition of $A_k(F; Y_i)$

$$\sum_{i=1}^{2n-1} |A_k(F; Y_i)| = 2 |A_k(F)| .$$

Assume that every $|A_k(F; Y_i)| \in P_k(F)$ is odd.

Then since $(2n-1)$ is odd $\sum_{i=1}^{2n-1} |A_k(F; Y_i)|$ must also be odd, which is a contradiction. Therefore at least one $|A_k(F; Y_i)| \in A_k(F)$ is even. ■

The following lemma enables us to subdivide the complete set of equivalence classes of 1-factorizations of K_{2n} into n subsets.

LEMMA 3.12. Let Φ represent the complete set of 1-factorizations of K_{2n} . Then Φ can be partitioned into n subsets as follows:

$\Phi(k,0) = \{F \in \Phi : \text{there exists a } F_j \in L(F) \text{ such that } |A_k(F_j; Y_1)| = 0\}.$

For $\mu = 1, \dots, n-1,$

$\Phi(k,\mu) = \{F \in \Phi : \text{there exists a } F_j \in L(F) \text{ such that } |A_k(F_j; Y_1)| = 2\mu \text{ and}$

$\frac{1}{2}|A_k(F_j; Y_i)| \notin \{0, \dots, \mu-1\} \text{ for all } i = 2, \dots, 2n-1\}.$

PROOF. It is clear from the definition of $\Phi(k,\mu)$ for $\mu = 0, \dots, n-1$ that if $F \in \Phi(k,\mu)$ then $L(F) \subseteq \Phi(k,\mu)$. By Lemma 3.11 we know that for every $F \in \Phi$ and every even k such that $4 \leq k \leq 2n$, some $|A_k(F; Y_1)| = 2\mu$ for some $\mu \in \{0, 1, \dots, n-1\}$. By Lemma 3.10 we know that if $|A_k(F; Y_1)| = 2\mu$ for some $\mu \in \{0, 1, \dots, n-1\}$ there exists a $F_j \in L(F)$ such that $|A_k(F_j; Y_1)| = 2\mu$. Clearly, therefore, $\Phi = \Phi(k,0) \cup \Phi(k,1) \cup \dots \cup \Phi(k,n-1)$. But, for $\Phi(k,\gamma)$ and $\Phi(k,\delta)$ such that $\gamma, \delta \in \{0, 1, \dots, n-1\}$ and $\delta > \gamma$ we have $2\gamma \in P_k(F_j)$ for every $F_j \in \Phi(k,\gamma)$ whereas $2\gamma \notin P_k(F_m)$ for every $F_m \in \Phi(k,\delta)$. Therefore, $\Phi(k,\gamma) \cap \Phi(k,\delta) = \emptyset$ for every $\gamma, \delta \in \{0, 1, \dots, n-1\}$ and $\gamma \neq \delta$. ■

THEOREM 3.13. There are 396 equivalence classes of 1-factorizations of K_{10} .

PROOF. In our procedure for enumerating equivalence classes of 1-factorizations of K_{10} , we considered alternating Hamiltonian circuits on the 1-factors and alternating 4-circuits on the vertices. We partitioned the set of equivalence classes into five subsets, $\Phi(10,0), \dots, \Phi(10,4)$, based on Lemma 3.12. Using Lemmas 3.6, 3.7 and 3.8 we only considered schedules which included the 1-factors listed in Figure 3.3. Computer Program 2 followed the backtracking technique described in Figure 3.1 but restricted itself to only produce schedules such that X_1 formed 2μ Hamiltonian circuits, where the value 2μ for $\mu \in \{0, \dots, 4\}$ was read into the program from a punched card.

As a schedule was generated by the computer program, it was tested to see if it represented a new equivalence class. Let F represent the current schedule being tested and let F' represent a previously-generated representative of an equivalence class. By Lemma 3.9 the program performed the following tests for isomorphism:

- 1) $|A_{10}(F)| = |A_{10}(F')|$.
- 2) $P_{10}(F) = P_{10}(F')$ and $Q_4(F) = Q_4(F')$.

If F and F' matched on these tests then the computer program performed a test based on part iii of Lemma 3.9. That is, it performed mappings on alternating circuits of F and F' in an order based on numbers of low multiplicity in $P_{10}(F)$ and $Q_4(F)$. If a mapping ψ was found such that $\psi F = F'$ then the count for the equivalence class was incremented by one. If no mapping could be found, the computer program created a new equivalence class with F as its representative. ■

See Appendix A for a listing of representatives of each equivalence class of 1-factorizations of K_{10} . Included with each schedule F are the sets $P_{10}(F)$ and $Q_4(F)$ as well as the number of Hamiltonian circuits, $|A_{10}(F)|$.

THEOREM 3.14. The order of the automorphism group of each of the 396 equivalence classes of 1-factorizations of K_{10} is as given in Appendix A.

PROOF. By Theorem 3.13, Computer Program 2 generated and counted all schedules, F , in each equivalence class which have the following characteristics:

- 1) $X_1 \in F$.
- 2) X_2 or X_3 is in F where X_2, X_3 are given in Figure 3.3.
- 3) If $X_2 \in F$ then one of $X_5, X_7, X_8, X_9, X_{10}, X_{11}$ or X_{12} listed in Figure 3.3 is in F . If $X_3 \in F$ then one of $X_4, \dots, X_7, X_{13}, \dots, X_{36}$ is in F .
- 4) For $\mu = 0, \dots, n-1$, $F \in \Phi(10, \mu)$ and $|A_{10}(F; X_1)| = 2\mu$. (See Lemma 3.12).

By Lemmas 3.6, 3.7, 3.8, each total generated by the computer program can be multiplied by the multiplication factors M_1, M_2 and M_3 or M_4 from Figure 3.3. By Lemma 3.10, each total can then be multiplied by 9 and divided by c (where c is the multiplicity of 2μ in $P_{10}(F)$) to give the total number of members in each equivalence class, $|L(F)|$.

Finally, by Theorem 2.16, $|G(F)| = \frac{10!}{|L(F)|}$. ■

THEOREM 3.15. There are 1,225,566,720 1-factorizations of K_{10} labelled on the set $\{1, 2, \dots, 10\}$.

PROOF. $\sum_{i=1}^{396} |L(F_i)| = 1,225,566,720$ where the 1-factorizations, F_i are listed in Appendix A and each $|L(F_i)|$ was found by Computer Program 2 by Theorem 3.14. ■

As noted earlier, listed with the representatives of the equivalence classes of 1-factorizations of K_{10} in Appendix A are the sets of numbers $P_{10}(F_i)$, (the set of counts of alternating Hamiltonian circuits across the 1-factors), and $Q_4(F_i)$, (the set of counts of alternating 4-circuits across the vertices). Unfortunately, these numbers do not completely characterize the equivalence classes. There are six pairs of equivalence classes with the same two sets of numbers (see Appendix A):

$\{F_{85}, F_{86}\}, \{F_{176}, F_{177}\}, \{F_{223}, F_{224}\}, \{F_{267}, F_{268}\},$
 $\{F_{290}, F_{291}\}, \{F_{303}, F_{304}\}$, and other tests must be performed in order to differentiate between them. We were unable to find an easily calculable method of completely characterizing the 396 equivalence classes of 1-factorizations of K_{10} .

It is of interest to note that the 1-factorization of K_{10} given by Theorem 2.6 is in the equivalence class $L(F_{290})$ where F_{290} is given in Appendix A. It has 27 Hamiltonian circuits and is therefore (as noted in Corollary 2.12) not an H-schedule. The order of its

automorphism group, $|G(\mathbb{F}_{290})|$, is 54.

We note also that K_{10} is the smallest complete graph which has a non-symmetrical 1-factorization (i.e., a 1-factorization which does not have any non-trivial automorphisms). In fact, of the 396 equivalence classes of 1-factorizations of K_{10} , 298 of them are non-symmetrical.

3-4. Conclusion

In our study we have counted the number and size of the equivalence classes of 1-factorizations of K_{2n} for $n \leq 5$. For larger n , the problem appears much more difficult, simply because of the large number of 1-factorizations involved. However, it should be possible to count the equivalence classes of 1-factorizations of K_{12} by a generalization of Lemmas 3.6, 3.7, and 3.8, and extending the method through each subset W_1 (see Sections 3-1 and 3-3 of Chapter 3) of the 1-factors of K_{12} . In other words, by constructively creating schedules round by round, testing for isomorphism at each stage, it should be possible to produce a complete set of non-equivalent schedules with 12 teams in a fairly "reasonable" length of time.

It seems clear, however, that the number of equivalence classes of 1-factorizations (as well as the total

number of 1-factorizations) of K_{2n} increases very rapidly as n increases. Therefore, to solve the problem for $2n$ larger than 12 by the methods of this thesis would appear to be extremely difficult, if not impossible.

CHAPTER 4

AN APPLICATION

Round robin schedules are commonly used for sporting matches and tournaments. Because of its ease of generation, the typical round robin schedule chosen is the scheme given by Theorem 2.6. We will show that, when given certain constraints, the choice of the round robin schedule can spell the difference between success or failure in solving the particular problem.

The problem is to create a schedule for sports such as curling, tennis, bowling, etc., where the number of playing fields are at a premium. The schedule must be as equitable as possible; that is, no team will have an advantage by playing at a particular favourable site more often than other teams. The schedule will be a round robin schedule involving $2n$ teams. Because each round will contain n games the number of fields to play on will also be n . In order to maximize the fairness of the schedule, the following two constraints are to be satisfied:

- (i) no team plays more than twice on the same field,
- (ii) no team plays twice in a row on the same field.

We chose the situation where we had 8 teams playing a round robin and 4 fields at our disposal. By Theorem 3.3

there are 6 equivalence classes of schedules with 8 teams (listed in Figure 3.2). A computer program was written to use a representative schedule from each of the six equivalence classes and to place games on fields round by round in an attempt to pass constraint (i). If it was successful it then reordered the rounds in an attempt to pass constraint (ii).

It was found that $L(F_2)$, $L(F_4)$ and $L(F_5)$ were the only round robin schedules which could pass the two constraints. The schedule $L(F_6)$ given by Theorem 2.6 passed constraint (i) but could not pass constraint (ii). The remaining two equivalence classes, $L(F_1)$ and $L(F_3)$, could not pass constraint (i).

Figure 4.1 below gives a round robin schedule with 8 teams and 4 fields which will pass the constraints. The schedule chosen is F_4 from Figure 3.2.

FIGURE 4.1 A ROUND ROBIN SCHEDULE WITH 8 TEAMS AND 4 FIELDS.

| <u>Round</u> | <u>Games</u> | <u>Field</u> | <u>Round</u> | <u>Games</u> | <u>Field</u> |
|--------------|--------------|--------------|--------------|--------------|--------------|
| 1 | 1 - 2 | 1 | 2 | 1 - 5 | 4 |
| | 3 - 4 | 2 | | 2 - 6 | 2 |
| | 5 - 6 | 3 | | 3 - 8 | 3 |
| | 7 - 8 | 4 | | 4 - 7 | 1 |

FIGURE 4.1 (cont'd)

| <u>Round</u> | <u>Games</u> | <u>Field</u> | <u>Round</u> | <u>Games</u> | <u>Field</u> |
|--------------|--------------|--------------|--------------|--------------|--------------|
| 3 | 1 - 3 | 1 | 4 | 1 - 4 | 2 |
| | 2 - 4 | 3 | | 2 - 3 | 4 |
| | 5 - 7 | 2 | | 5 - 8 | 1 |
| | 6 - 8 | 4 | | 6 - 7 | 3 |
| 5 | 1 - 7 | 4 | 6 | 1 - 8 | 3 |
| | 2 - 8 | 2 | | 2 - 5 | 1 |
| | 3 - 6 | 1 | | 3 - 7 | 2 |
| | 4 - 5 | 3 | | 4 - 6 | 4 |
| 7 | 1 - 6 | 2 | | | |
| | 2 - 7 | 3 | | | |
| | 3 - 5 | 4 | | | |
| | 4 - 8 | 1 | | | |

Note: Game $x - y$ corresponds to team x playing team y .

We note that the schedule is cyclic with respect to the rounds; i.e., the rounds can be played in the order $i \pmod{7}$, $i+1 \pmod{7}$, ..., $i+6 \pmod{7}$. Clearly they can also be played in the reverse order.

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APPENDIX A

Representatives from each of the equivalence classes of
1-factorizations of K_{10} .

Appendix A consists of 3 listings:

1. Factor list - a listing of the 1-factors X_i ($i = 1, \dots, 270$) which are contained in the subsequent schedules.
2. Schedules by factor number - a listing of the 396 representative schedules. The numbers under the columns labelled "FACTORS" refer to the 1-factors in the Factor list. The "ORDER" is the order of the automorphism group. The column "HC" gives the total number of alternating Hamiltonian circuits in the schedule. The columns labelled "# HAMILT CIRCUITS ACROSS THE FCTRS" gives the numbers in $P_{10}(F_i)$ ordered by factor number. The columns labelled "# 4-CIRCUITS ACROSS THE VERTICES" give the numbers in $Q_4(F_i)$ ordered by vertex number.
3. Auxiliary listing to the list of schedules - The numbers in $P_{10}(F_i)$ and $Q_4(F_i)$ have been ordered from smallest to largest. This listing gives a partial characterization of the 396 equivalence classes.

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
 FACTOR LIST

| FCTR | EDGES | FCTR | EDGES | FCTR | EDGES |
|------|-------|------|-------|------|-------|
| 51 | 1-6 | 92 | 1-7 | 5-7 | 7-10 |
| 53 | 1-7 | 94 | 1-7 | 6-9 | 8-9 |
| 55 | 1-7 | 96 | 1-7 | 6-10 | 6-10 |
| 57 | 1-7 | 98 | 1-7 | 5-8 | 5-10 |
| 59 | 1-7 | 100 | 1-7 | 5-9 | 6-8 |
| 101 | 1-7 | 102 | 1-7 | 6-10 | 6-10 |
| 103 | 1-7 | 104 | 1-7 | 8-9 | 8-9 |
| 105 | 1-7 | 106 | 1-7 | 8-10 | 6-10 |
| 107 | 1-7 | 108 | 1-7 | 9-9 | 6-10 |
| 109 | 1-7 | 110 | 1-7 | 9-10 | 6-10 |
| 111 | 1-7 | 112 | 1-7 | 5-8 | 5-10 |
| 113 | 1-7 | 114 | 1-7 | 5-9 | 5-10 |
| 115 | 1-7 | 116 | 1-7 | 6-10 | 5-10 |
| 117 | 1-7 | 118 | 1-7 | 6-10 | 5-10 |
| 119 | 1-7 | 120 | 1-7 | 6-10 | 5-10 |
| 121 | 1-7 | 122 | 1-7 | 6-10 | 5-10 |
| 123 | 1-7 | 124 | 1-7 | 6-10 | 5-10 |
| 125 | 1-7 | 126 | 1-7 | 6-10 | 5-10 |
| 127 | 1-7 | 128 | 1-7 | 6-10 | 5-10 |
| 129 | 1-7 | 130 | 1-7 | 6-10 | 5-10 |
| 131 | 1-7 | 132 | 1-7 | 6-10 | 5-10 |
| 133 | 1-8 | 134 | 1-8 | 6-10 | 5-10 |
| 135 | 1-8 | 136 | 1-8 | 6-10 | 5-10 |
| 137 | 1-8 | 138 | 1-8 | 6-10 | 5-10 |
| 139 | 1-8 | 140 | 1-8 | 6-10 | 5-10 |
| 141 | 1-8 | 142 | 1-8 | 6-10 | 5-10 |
| 143 | 1-8 | 144 | 1-8 | 6-10 | 5-10 |
| 145 | 1-8 | 146 | 1-8 | 6-10 | 5-10 |
| 147 | 1-8 | 148 | 1-8 | 6-10 | 5-10 |
| 149 | 1-8 | 150 | 1-8 | 6-10 | 5-10 |
| 151 | 1-8 | 152 | 1-8 | 6-10 | 5-10 |
| 153 | 1-8 | 154 | 1-8 | 6-10 | 5-10 |
| 155 | 1-8 | 156 | 1-8 | 6-10 | 5-10 |
| 157 | 1-8 | 158 | 1-8 | 6-10 | 5-10 |
| 159 | 1-8 | 160 | 1-8 | 6-10 | 5-10 |
| 161 | 1-8 | 162 | 1-8 | 6-10 | 5-10 |
| 163 | 1-8 | 164 | 1-8 | 6-10 | 5-10 |
| 165 | 1-8 | 166 | 1-8 | 6-10 | 5-10 |
| 167 | 1-8 | 168 | 1-8 | 6-10 | 5-10 |
| 169 | 1-8 | 170 | 1-8 | 6-10 | 5-10 |
| 171 | 1-8 | 172 | 1-8 | 6-10 | 5-10 |
| 173 | 1-8 | 174 | 1-8 | 6-10 | 5-10 |
| 175 | 1-8 | 176 | 1-8 | 6-10 | 5-10 |
| 177 | 1-8 | 178 | 1-8 | 6-10 | 5-10 |
| 179 | 1-8 | 180 | 1-8 | 6-10 | 5-10 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10

FACTOR LIST

| FCTR | EDGES | FCTR | EDGES |
|------|-------|------|-------|
| 181 | 1-1-9 | 182 | 1-1-9 |
| 183 | 1-1-9 | 184 | 1-1-9 |
| 185 | 1-1-9 | 186 | 1-1-9 |
| 187 | 1-1-9 | 188 | 1-1-9 |
| 189 | 1-1-9 | 190 | 1-1-9 |
| 191 | 1-1-9 | 192 | 1-1-9 |
| 193 | 1-1-9 | 194 | 1-1-9 |
| 195 | 1-1-9 | 196 | 1-1-9 |
| 197 | 1-1-9 | 198 | 1-1-9 |
| 199 | 1-1-9 | 200 | 1-1-9 |
| 201 | 1-1-9 | 202 | 1-1-9 |
| 203 | 1-1-9 | 204 | 1-1-9 |
| 205 | 1-1-9 | 206 | 1-1-9 |
| 207 | 1-1-9 | 208 | 1-1-9 |
| 209 | 1-1-9 | 210 | 1-1-9 |
| 211 | 1-1-9 | 212 | 1-1-9 |
| 213 | 1-1-9 | 214 | 1-1-9 |
| 215 | 1-1-9 | 216 | 1-1-9 |
| 217 | 1-1-9 | 218 | 1-1-9 |
| 219 | 1-1-9 | 220 | 1-1-9 |
| 221 | 1-1-9 | 222 | 1-1-9 |
| 223 | 1-1-9 | 224 | 1-1-9 |
| 225 | 1-1-9 | 226 | 1-1-9 |
| 227 | 1-1-9 | 228 | 1-1-9 |
| 229 | 1-1-9 | 230 | 1-1-9 |
| 231 | 1-1-9 | 232 | 1-1-9 |
| 233 | 1-1-9 | 234 | 1-1-9 |
| 235 | 1-1-9 | 236 | 1-1-9 |
| 237 | 1-1-9 | 238 | 1-1-9 |
| 239 | 1-1-9 | 240 | 1-1-9 |
| 241 | 1-1-9 | 242 | 1-1-9 |
| 243 | 1-1-9 | 244 | 1-1-9 |
| 245 | 1-1-9 | 246 | 1-1-9 |
| 247 | 1-1-9 | 248 | 1-1-9 |
| 249 | 1-1-9 | 250 | 1-1-9 |
| 251 | 1-1-9 | 252 | 1-1-9 |
| 253 | 1-1-9 | 254 | 1-1-9 |
| 255 | 1-1-9 | 256 | 1-1-9 |
| 257 | 1-1-9 | 258 | 1-1-9 |
| 259 | 1-1-9 | 260 | 1-1-9 |
| 261 | 1-1-9 | 262 | 1-1-9 |
| 263 | 1-1-9 | 264 | 1-1-9 |
| 265 | 1-1-9 | 266 | 1-1-9 |
| 267 | 1-1-9 | 268 | 1-1-9 |
| 269 | 1-1-9 | 270 | 1-1-9 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES BY FACTOR NUMBER

| SCHED | ORDER | MC | # FAMILY CIRCUITS ACROSS THE FCTRS | # 4-CIRCUITS ACROSS THE VERTICES |
|-------|-------|----|------------------------------------|----------------------------------|
| 137 | 11 | 24 | 44 | 77 |
| 138 | 11 | 24 | 44 | 77 |
| 139 | 11 | 24 | 44 | 77 |
| 140 | 11 | 24 | 44 | 77 |
| 141 | 11 | 24 | 44 | 77 |
| 142 | 11 | 24 | 44 | 77 |
| 143 | 11 | 24 | 44 | 77 |
| 144 | 11 | 24 | 44 | 77 |
| 145 | 11 | 24 | 44 | 77 |
| 146 | 11 | 24 | 44 | 77 |
| 147 | 11 | 24 | 44 | 77 |
| 148 | 11 | 24 | 44 | 77 |
| 149 | 11 | 24 | 44 | 77 |
| 150 | 11 | 24 | 44 | 77 |
| 151 | 11 | 24 | 44 | 77 |
| 152 | 11 | 24 | 44 | 77 |
| 153 | 11 | 24 | 44 | 77 |
| 154 | 11 | 24 | 44 | 77 |
| 155 | 11 | 24 | 44 | 77 |
| 156 | 11 | 24 | 44 | 77 |
| 157 | 11 | 24 | 44 | 77 |
| 158 | 11 | 24 | 44 | 77 |
| 159 | 11 | 24 | 44 | 77 |
| 160 | 11 | 24 | 44 | 77 |
| 161 | 11 | 24 | 44 | 77 |
| 162 | 11 | 24 | 44 | 77 |
| 163 | 11 | 24 | 44 | 77 |
| 164 | 11 | 24 | 44 | 77 |
| 165 | 11 | 24 | 44 | 77 |
| 166 | 11 | 24 | 44 | 77 |
| 167 | 11 | 24 | 44 | 77 |
| 168 | 11 | 24 | 44 | 77 |
| 169 | 11 | 24 | 44 | 77 |
| 170 | 11 | 24 | 44 | 77 |
| 171 | 11 | 24 | 44 | 77 |
| 172 | 11 | 24 | 44 | 77 |
| 173 | 11 | 24 | 44 | 77 |
| 174 | 11 | 24 | 44 | 77 |
| 175 | 11 | 24 | 44 | 77 |
| 176 | 11 | 24 | 44 | 77 |
| 177 | 11 | 24 | 44 | 77 |
| 178 | 11 | 24 | 44 | 77 |
| 179 | 11 | 24 | 44 | 77 |
| 180 | 11 | 24 | 44 | 77 |
| 181 | 11 | 24 | 44 | 77 |
| 182 | 11 | 24 | 44 | 77 |
| 183 | 11 | 24 | 44 | 77 |
| 184 | 11 | 24 | 44 | 77 |
| 185 | 11 | 24 | 44 | 77 |
| 186 | 11 | 24 | 44 | 77 |
| 187 | 11 | 24 | 44 | 77 |
| 188 | 11 | 24 | 44 | 77 |
| 189 | 11 | 24 | 44 | 77 |
| 190 | 11 | 24 | 44 | 77 |
| 191 | 11 | 24 | 44 | 77 |
| 192 | 11 | 24 | 44 | 77 |
| 193 | 11 | 24 | 44 | 77 |
| 194 | 11 | 24 | 44 | 77 |
| 195 | 11 | 24 | 44 | 77 |
| 196 | 11 | 24 | 44 | 77 |
| 197 | 11 | 24 | 44 | 77 |
| 198 | 11 | 24 | 44 | 77 |
| 199 | 11 | 24 | 44 | 77 |
| 200 | 11 | 24 | 44 | 77 |
| 201 | 11 | 24 | 44 | 77 |
| 202 | 11 | 24 | 44 | 77 |
| 203 | 11 | 24 | 44 | 77 |
| 204 | 11 | 24 | 44 | 77 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES BY FACTOR NUMBER

| SCHED | * | * | * | * | * | FACTORS | * | * | ORDER | PC | # HAMILT CIRCUITS ACROSS THE FCTRS | # HAMILT CIRCUITS ACROSS THE VERTICES | # 4-CIRCUITS ACROSS THE VERTICES |
|-------|---|---|----|---|---|---------------------|---|---|-------|----|------------------------------------|---------------------------------------|----------------------------------|
| 273 | 1 | 2 | 1 | 1 | 1 | 65 119 177 182 242 | 1 | 1 | 1 | 27 | 4 | 5 | |
| 274 | 1 | 3 | 10 | 1 | 1 | 79 121 142 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 275 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 276 | 1 | 3 | 10 | 1 | 1 | 126 169 184 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 277 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 278 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 279 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 280 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 281 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 282 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 283 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 284 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 285 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 286 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 287 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 288 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 289 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 290 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 291 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 292 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 293 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 294 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 295 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 296 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 297 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 298 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 299 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 300 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 301 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 302 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 303 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 304 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 305 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 306 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 307 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 308 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 309 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 310 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 311 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 312 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 313 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 314 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 315 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 316 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 317 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 318 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 319 | 1 | 3 | 10 | 1 | 1 | 177 187 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |
| 320 | 1 | 3 | 10 | 1 | 1 | 149 184 201 222 227 | 1 | 1 | 1 | 27 | 7 | 5 | |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES BY FACTOR NUMBER

| SCHED | * | * | * | * | FACTORS | * * * | ORDER | FC | # HAMILT CIRCUITS ACROSS THE FCTRS | # 4-CIRCUITS ACROSS THE VERTICES |
|-------|---|---|---|---|---------|-------|-------|----|------------------------------------|----------------------------------|
| 141 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 142 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 143 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 144 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 145 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 146 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 147 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 148 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 149 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 150 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 151 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 152 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 153 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 154 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 155 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 156 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 157 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 158 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 159 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 160 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 161 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 162 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 163 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 164 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 165 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 166 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 167 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 168 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 169 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 170 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 171 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 172 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 173 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 174 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 175 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 176 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 177 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 178 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 179 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 180 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 181 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 182 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 183 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 184 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 185 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 186 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 187 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 188 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 189 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 190 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 191 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 192 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 193 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 194 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 195 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 196 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 197 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 198 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 199 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |
| 200 | 1 | 1 | 1 | 1 | 59 | 176 | 1 | 29 | 4 | 3 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES WITH CIRCUITS REORDERED FOR EASIER REFERENCE - REFER TO PREVIOUS LISTING

| SCHED | HC | # HAMILT CIRCUITS ACROSS THE FCTRS | # 4-CIRCUITS ACROSS THE VERTICES | 36 |
|-------|-----|------------------------------------|----------------------------------|----|
| 1 | 2 | 0 | 12 | 12 |
| 2 | 3 | 0 | 12 | 12 |
| 3 | 4 | 0 | 12 | 12 |
| 4 | 5 | 0 | 12 | 12 |
| 5 | 6 | 0 | 12 | 12 |
| 6 | 7 | 0 | 12 | 12 |
| 7 | 8 | 0 | 12 | 12 |
| 8 | 9 | 0 | 12 | 12 |
| 9 | 10 | 0 | 12 | 12 |
| 10 | 11 | 0 | 12 | 12 |
| 11 | 12 | 0 | 12 | 12 |
| 12 | 13 | 0 | 12 | 12 |
| 13 | 14 | 0 | 12 | 12 |
| 14 | 15 | 0 | 12 | 12 |
| 15 | 16 | 0 | 12 | 12 |
| 16 | 17 | 0 | 12 | 12 |
| 17 | 18 | 0 | 12 | 12 |
| 18 | 19 | 0 | 12 | 12 |
| 19 | 20 | 0 | 12 | 12 |
| 20 | 21 | 0 | 12 | 12 |
| 21 | 22 | 0 | 12 | 12 |
| 22 | 23 | 0 | 12 | 12 |
| 23 | 24 | 0 | 12 | 12 |
| 24 | 25 | 0 | 12 | 12 |
| 25 | 26 | 0 | 12 | 12 |
| 26 | 27 | 0 | 12 | 12 |
| 27 | 28 | 0 | 12 | 12 |
| 28 | 29 | 0 | 12 | 12 |
| 29 | 30 | 0 | 12 | 12 |
| 30 | 31 | 0 | 12 | 12 |
| 31 | 32 | 0 | 12 | 12 |
| 32 | 33 | 0 | 12 | 12 |
| 33 | 34 | 0 | 12 | 12 |
| 34 | 35 | 0 | 12 | 12 |
| 35 | 36 | 0 | 12 | 12 |
| 36 | 37 | 0 | 12 | 12 |
| 37 | 38 | 0 | 12 | 12 |
| 38 | 39 | 0 | 12 | 12 |
| 39 | 40 | 0 | 12 | 12 |
| 40 | 41 | 0 | 12 | 12 |
| 41 | 42 | 0 | 12 | 12 |
| 42 | 43 | 0 | 12 | 12 |
| 43 | 44 | 0 | 12 | 12 |
| 44 | 45 | 0 | 12 | 12 |
| 45 | 46 | 0 | 12 | 12 |
| 46 | 47 | 0 | 12 | 12 |
| 47 | 48 | 0 | 12 | 12 |
| 48 | 49 | 0 | 12 | 12 |
| 49 | 50 | 0 | 12 | 12 |
| 50 | 51 | 0 | 12 | 12 |
| 51 | 52 | 0 | 12 | 12 |
| 52 | 53 | 0 | 12 | 12 |
| 53 | 54 | 0 | 12 | 12 |
| 54 | 55 | 0 | 12 | 12 |
| 55 | 56 | 0 | 12 | 12 |
| 56 | 57 | 0 | 12 | 12 |
| 57 | 58 | 0 | 12 | 12 |
| 58 | 59 | 0 | 12 | 12 |
| 59 | 60 | 0 | 12 | 12 |
| 60 | 61 | 0 | 12 | 12 |
| 61 | 62 | 0 | 12 | 12 |
| 62 | 63 | 0 | 12 | 12 |
| 63 | 64 | 0 | 12 | 12 |
| 64 | 65 | 0 | 12 | 12 |
| 65 | 66 | 0 | 12 | 12 |
| 66 | 67 | 0 | 12 | 12 |
| 67 | 68 | 0 | 12 | 12 |
| 68 | 69 | 0 | 12 | 12 |
| 69 | 70 | 0 | 12 | 12 |
| 70 | 71 | 0 | 12 | 12 |
| 71 | 72 | 0 | 12 | 12 |
| 72 | 73 | 0 | 12 | 12 |
| 73 | 74 | 0 | 12 | 12 |
| 74 | 75 | 0 | 12 | 12 |
| 75 | 76 | 0 | 12 | 12 |
| 76 | 77 | 0 | 12 | 12 |
| 77 | 78 | 0 | 12 | 12 |
| 78 | 79 | 0 | 12 | 12 |
| 79 | 80 | 0 | 12 | 12 |
| 80 | 81 | 0 | 12 | 12 |
| 81 | 82 | 0 | 12 | 12 |
| 82 | 83 | 0 | 12 | 12 |
| 83 | 84 | 0 | 12 | 12 |
| 84 | 85 | 0 | 12 | 12 |
| 85 | 86 | 0 | 12 | 12 |
| 86 | 87 | 0 | 12 | 12 |
| 87 | 88 | 0 | 12 | 12 |
| 88 | 89 | 0 | 12 | 12 |
| 89 | 90 | 0 | 12 | 12 |
| 90 | 91 | 0 | 12 | 12 |
| 91 | 92 | 0 | 12 | 12 |
| 92 | 93 | 0 | 12 | 12 |
| 93 | 94 | 0 | 12 | 12 |
| 94 | 95 | 0 | 12 | 12 |
| 95 | 96 | 0 | 12 | 12 |
| 96 | 97 | 0 | 12 | 12 |
| 97 | 98 | 0 | 12 | 12 |
| 98 | 99 | 0 | 12 | 12 |
| 99 | 100 | 0 | 12 | 12 |
| 100 | 101 | 0 | 12 | 12 |
| 101 | 102 | 0 | 12 | 12 |
| 102 | 103 | 0 | 12 | 12 |
| 103 | 104 | 0 | 12 | 12 |
| 104 | 105 | 0 | 12 | 12 |
| 105 | 106 | 0 | 12 | 12 |
| 106 | 107 | 0 | 12 | 12 |
| 107 | 108 | 0 | 12 | 12 |
| 108 | 109 | 0 | 12 | 12 |
| 109 | 110 | 0 | 12 | 12 |
| 110 | 111 | 0 | 12 | 12 |
| 111 | 112 | 0 | 12 | 12 |
| 112 | 113 | 0 | 12 | 12 |
| 113 | 114 | 0 | 12 | 12 |
| 114 | 115 | 0 | 12 | 12 |
| 115 | 116 | 0 | 12 | 12 |
| 116 | 117 | 0 | 12 | 12 |
| 117 | 118 | 0 | 12 | 12 |
| 118 | 119 | 0 | 12 | 12 |
| 119 | 120 | 0 | 12 | 12 |
| 120 | 121 | 0 | 12 | 12 |
| 121 | 122 | 0 | 12 | 12 |
| 122 | 123 | 0 | 12 | 12 |
| 123 | 124 | 0 | 12 | 12 |
| 124 | 125 | 0 | 12 | 12 |
| 125 | 126 | 0 | 12 | 12 |
| 126 | 127 | 0 | 12 | 12 |
| 127 | 128 | 0 | 12 | 12 |
| 128 | 129 | 0 | 12 | 12 |
| 129 | 130 | 0 | 12 | 12 |
| 130 | 131 | 0 | 12 | 12 |
| 131 | 132 | 0 | 12 | 12 |
| 132 | 133 | 0 | 12 | 12 |
| 133 | 134 | 0 | 12 | 12 |
| 134 | 135 | 0 | 12 | 12 |
| 135 | 136 | 0 | 12 | 12 |
| 136 | 137 | 0 | 12 | 12 |
| 137 | 138 | 0 | 12 | 12 |
| 138 | 139 | 0 | 12 | 12 |
| 139 | 140 | 0 | 12 | 12 |
| 140 | 141 | 0 | 12 | 12 |
| 141 | 142 | 0 | 12 | 12 |
| 142 | 143 | 0 | 12 | 12 |
| 143 | 144 | 0 | 12 | 12 |
| 144 | 145 | 0 | 12 | 12 |
| 145 | 146 | 0 | 12 | 12 |
| 146 | 147 | 0 | 12 | 12 |
| 147 | 148 | 0 | 12 | 12 |
| 148 | 149 | 0 | 12 | 12 |
| 149 | 150 | 0 | 12 | 12 |
| 150 | 151 | 0 | 12 | 12 |
| 151 | 152 | 0 | 12 | 12 |
| 152 | 153 | 0 | 12 | 12 |
| 153 | 154 | 0 | 12 | 12 |
| 154 | 155 | 0 | 12 | 12 |
| 155 | 156 | 0 | 12 | 12 |
| 156 | 157 | 0 | 12 | 12 |
| 157 | 158 | 0 | 12 | 12 |
| 158 | 159 | 0 | 12 | 12 |
| 159 | 160 | 0 | 12 | 12 |
| 160 | 161 | 0 | 12 | 12 |
| 161 | 162 | 0 | 12 | 12 |
| 162 | 163 | 0 | 12 | 12 |
| 163 | 164 | 0 | 12 | 12 |
| 164 | 165 | 0 | 12 | 12 |
| 165 | 166 | 0 | 12 | 12 |
| 166 | 167 | 0 | 12 | 12 |
| 167 | 168 | 0 | 12 | 12 |
| 168 | 169 | 0 | 12 | 12 |
| 169 | 170 | 0 | 12 | 12 |
| 170 | 171 | 0 | 12 | 12 |
| 171 | 172 | 0 | 12 | 12 |
| 172 | 173 | 0 | 12 | 12 |
| 173 | 174 | 0 | 12 | 12 |
| 174 | 175 | 0 | 12 | 12 |
| 175 | 176 | 0 | 12 | 12 |
| 176 | 177 | 0 | 12 | 12 |
| 177 | 178 | 0 | 12 | 12 |
| 178 | 179 | 0 | 12 | 12 |
| 179 | 180 | 0 | 12 | 12 |
| 180 | 181 | 0 | 12 | 12 |
| 181 | 182 | 0 | 12 | 12 |
| 182 | 183 | 0 | 12 | 12 |
| 183 | 184 | 0 | 12 | 12 |
| 184 | 185 | 0 | 12 | 12 |
| 185 | 186 | 0 | 12 | 12 |
| 186 | 187 | 0 | 12 | 12 |
| 187 | 188 | 0 | 12 | 12 |
| 188 | 189 | 0 | 12 | 12 |
| 189 | 190 | 0 | 12 | 12 |
| 190 | 191 | 0 | 12 | 12 |
| 191 | 192 | 0 | 12 | 12 |
| 192 | 193 | 0 | 12 | 12 |
| 193 | 194 | 0 | 12 | 12 |
| 194 | 195 | 0 | 12 | 12 |
| 195 | 196 | 0 | 12 | 12 |
| 196 | 197 | 0 | 12 | 12 |
| 197 | 198 | 0 | 12 | 12 |
| 198 | 199 | 0 | 12 | 12 |
| 199 | 200 | 0 | 12 | 12 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES WITH CIRCUITS REORDERED FOR EASIER REFERENCE - REFER TO PREVIOUS LISTING
4-CIRCUITS ACROSS THE VERTICES

| SCHED | HC | # | FAMILT | CIRCUITS | ACROSS | THE | FCTRS | # | 4-CIRCUITS | ACROSS | THE | VERTICES |
|-------|----|---|--------|----------|--------|-----|-------|---|------------|--------|-----|----------|
| 205 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 206 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 207 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 208 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 209 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 210 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 211 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 212 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 213 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 214 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 215 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 216 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 217 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 218 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 219 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 220 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 221 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 222 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 223 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 224 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 225 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 226 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 227 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 228 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 229 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 230 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 231 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 232 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 233 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 234 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 235 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 236 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 237 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 238 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 239 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 240 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 241 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 242 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 243 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 244 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 245 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 246 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 247 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 248 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 249 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 250 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 251 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 252 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 253 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 254 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 255 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 256 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 257 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 258 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 259 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 260 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 261 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 262 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 263 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 264 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 265 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 266 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 267 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 268 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 269 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |
| 270 | 26 | 4 | 4 | 5 | 5 | 6 | 6 | 7 | 7 | 7 | 7 | 7 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES WITH CIRCUITS REORDERED FOR EASIER REFERENCE - REFER TO PREVIOUS LISTING

| SCHED | HC | # HAMILT CIRCUITS ACROSS THE FCTRS | # 4-CIRCUITS ACROSS THE VERTICES |
|-------|----|------------------------------------|----------------------------------|
| 273 | 27 | 4 | 5 |
| 274 | 27 | 5 | 5 |
| 275 | 27 | 5 | 5 |
| 276 | 27 | 5 | 5 |
| 277 | 27 | 5 | 5 |
| 278 | 27 | 5 | 5 |
| 279 | 27 | 5 | 5 |
| 280 | 27 | 5 | 5 |
| 281 | 27 | 5 | 5 |
| 282 | 27 | 5 | 5 |
| 283 | 27 | 5 | 5 |
| 284 | 27 | 5 | 5 |
| 285 | 27 | 5 | 5 |
| 286 | 27 | 5 | 5 |
| 287 | 27 | 5 | 5 |
| 288 | 27 | 5 | 5 |
| 289 | 27 | 5 | 5 |
| 290 | 27 | 5 | 5 |
| 291 | 28 | 4 | 4 |
| 292 | 28 | 4 | 4 |
| 293 | 28 | 4 | 4 |
| 294 | 28 | 4 | 4 |
| 295 | 28 | 4 | 4 |
| 296 | 28 | 4 | 4 |
| 297 | 28 | 4 | 4 |
| 298 | 28 | 4 | 4 |
| 299 | 28 | 4 | 4 |
| 300 | 28 | 4 | 4 |
| 301 | 28 | 4 | 4 |
| 302 | 28 | 4 | 4 |
| 303 | 28 | 4 | 4 |
| 304 | 28 | 4 | 4 |
| 305 | 28 | 4 | 4 |
| 306 | 28 | 4 | 4 |
| 307 | 28 | 4 | 4 |
| 308 | 28 | 4 | 4 |
| 309 | 28 | 4 | 4 |
| 310 | 28 | 4 | 4 |
| 311 | 28 | 4 | 4 |
| 312 | 28 | 4 | 4 |
| 313 | 28 | 4 | 4 |
| 314 | 28 | 4 | 4 |
| 315 | 28 | 4 | 4 |
| 316 | 28 | 4 | 4 |
| 317 | 28 | 4 | 4 |
| 318 | 28 | 4 | 4 |
| 319 | 28 | 4 | 4 |
| 320 | 28 | 4 | 4 |
| 321 | 28 | 4 | 4 |
| 322 | 28 | 4 | 4 |
| 323 | 28 | 4 | 4 |
| 324 | 28 | 4 | 4 |
| 325 | 28 | 4 | 4 |
| 326 | 28 | 4 | 4 |
| 327 | 28 | 4 | 4 |
| 328 | 28 | 4 | 4 |
| 329 | 28 | 4 | 4 |
| 330 | 28 | 4 | 4 |
| 331 | 28 | 4 | 4 |
| 332 | 28 | 4 | 4 |
| 333 | 28 | 4 | 4 |
| 334 | 28 | 4 | 4 |
| 335 | 28 | 4 | 4 |
| 336 | 28 | 4 | 4 |
| 337 | 28 | 4 | 4 |
| 338 | 28 | 4 | 4 |
| 339 | 28 | 4 | 4 |
| 340 | 28 | 4 | 4 |

REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR SCHEDULES OF ORDER 10
SCHEDULES WITH CIRCUITS REQUIRED FOR EASIR REFERENCE - REFER TO PREVIOUS LISTING

| SCHED HC | # HAMILT CIRCUITS ACROSS THE FCTRS | # 4-CIRCUITS ACROSS THE VERTICES |
|----------|------------------------------------|----------------------------------|
| 241 | 4 | 4 |
| 242 | 5 | 4 |
| 243 | 5 | 5 |
| 244 | 5 | 4 |
| 245 | 5 | 5 |
| 246 | 5 | 5 |
| 247 | 5 | 5 |
| 248 | 5 | 5 |
| 249 | 5 | 5 |
| 250 | 5 | 5 |
| 251 | 5 | 5 |
| 252 | 5 | 5 |
| 253 | 5 | 5 |
| 254 | 5 | 5 |
| 255 | 5 | 5 |
| 256 | 5 | 5 |
| 257 | 5 | 5 |
| 258 | 5 | 5 |
| 259 | 5 | 5 |
| 260 | 5 | 5 |
| 261 | 5 | 5 |
| 262 | 5 | 5 |
| 263 | 5 | 5 |
| 264 | 5 | 5 |
| 265 | 5 | 5 |
| 266 | 5 | 5 |
| 267 | 5 | 5 |
| 268 | 5 | 5 |
| 269 | 5 | 5 |
| 270 | 5 | 5 |
| 271 | 5 | 5 |
| 272 | 5 | 5 |
| 273 | 5 | 5 |
| 274 | 5 | 5 |
| 275 | 5 | 5 |
| 276 | 5 | 5 |
| 277 | 5 | 5 |
| 278 | 5 | 5 |
| 279 | 5 | 5 |
| 280 | 5 | 5 |
| 281 | 5 | 5 |
| 282 | 5 | 5 |
| 283 | 5 | 5 |
| 284 | 5 | 5 |
| 285 | 5 | 5 |
| 286 | 5 | 5 |
| 287 | 5 | 5 |
| 288 | 5 | 5 |
| 289 | 5 | 5 |
| 290 | 5 | 5 |
| 291 | 5 | 5 |
| 292 | 5 | 5 |
| 293 | 5 | 5 |
| 294 | 5 | 5 |
| 295 | 5 | 5 |
| 296 | 5 | 5 |

APPENDIX B

Computer Program 1

The program (SCHPGM1) finds and lists all schedules with $2n$ teams which include the fixed 1-factor X_1 . It tests all the schedules for isomorphism and lists the members in each equivalence class.

The program reads input cards with the following information:

1. N = number of vertices (corresponds to $2n$ in the thesis body).
2. MAXSCH = maximum number of schedules to be generated and stored.

The program generates and lists all 1-factors in the routine FØRMFAC and stores them in the array FACTØR. It then forms schedules from these 1-factors in the routine FØRMSCH. The schedules are stored in the array SCHED. The routine TRANFAC aids in the equivalence testing of the schedules. Its purpose is to form an array (FTRANS) of all the mappings which can take place on the 1-factors. The actual isomorphism tests are done in the routine EQUIV. It calls two subroutines, FINDSCH and PRINTEQ, to find the equivalent schedules in the array SCHED and to print the listings of the equivalence classes.

The program was run on the 512K IBM 370-145 at the University of Victoria. It was compiled using the ØS/360 PL/1 F compiler. Compile time was 34.8 seconds. Core storage was allocated dynamically according to the values of N and MAXSCH read in. For N = 8 and MAXSCH = 500 the storage used was 128K and execution time was 2 minutes.

```

SCHPGM1: PROCEDURE OPTICNS (MAIN);

/* THIS PGM WILL GENERATE ALL 1-FACTORS FOR THE COMPLETE GRAPH WITH
N VERTICES. IT THEN FORMS ALL ROUND ROBIN SCHEDULES WHICH INCLUDE
FACTOR 1. IT ALSO FORMS EQUIVALENCE CLASSES AND PRINTS THEM OUT. */

DECLARE FACTOR (NFAC,N) FIXED BINARY CONTROLLED;
DECLARE FTRANS (NFAC,NEDGE) FIXED BINARY CONTROLLED;
DECLARE NSCH (NM2) FIXED BINARY CONTROLLED;
DECLARE SCHED (MAXSCH,N) FIXED BINARY CONTROLLED;

GET LIST (N);
GET LIST (MAXSCH);
NM1 = N-1;
NM2 = N-2;
NEF = 1;
NC = 1;
SS1: NC = NC + 2;
IF NC > N-2 THEN GO TO SS2;
NEF = NEF*NC;
GO TO SS1;
SS2: NFAC = (N-1)*NEF;
NEDGE = N*(N-1)/2;
NSS = 1;
NSF = 500;
ND2 = N/2;

CALL FORMFAC;
PUT SKIP(1) EDIT ('FACTOR','EDGES') (A,X(12),A);
PUT SKIP(1);
DO K = 1 TO NFAC;
  PUT SKIP(1) EDIT (K,(FACTOR(K,J) DO J = 1 TO N))
  (F(3),9 (X(3),F(3),F(3)));
END;

SS3: PUT PAGE EDIT ('LISTING OF SCHEDULES AND FACTOR NUMBERS FOR ORDER',N)
(A,F(3));
PUT SKIP(1);
CALL FORMSCH;
DO I = 1 TO NCNT;
  PUT SKIP(1) EDIT (I,'1',(SCHED(I,I1) DO I1 = 1 TO NM2))
  (F(5),X(10),A,17 F(5));
END;

CALL TRANFAC;

FREE FACTOR;

IF N>8 THEN GO TO SS4;
PUT PAGE EDIT ('FACTOR','TRANSFORMATIONS') (A,X(10),A);
PUT SKIP(2) EDIT (((I,'-',J DO J = I+1 TO N) DO I = 1 TO NM1))
(X(4),28 (F(1),A,F(1),X(1)));
PUT SKIP(1);
DO K = 1 TO NFAC;
  PUT SKIP(1) EDIT (K,(FTRANS(K,IJ) DO IJ = 1 TO NEDGE))
  (F(3),28 (X(1),F(3)));
END;

SS4: CALL EQUIV;

PUT SKIP (1);

PUT SKIP(1) EDIT ((I2,SCHED(I2,NM1) DO I2 = 1 TO NCNT))
(COL(1),10 (F(4),F(5),X(3)));

/* ***** */

FORMFAC: PROCEDURE;

/* TO GENERATE ALL THE FACTORS FOR N TEAMS */

DECLARE TAB(NTAB) FIXED BINARY CONTROLLED;
DECLARE (TE(ND2),TE(ND2),TP(ND2)) FIXED BINARY CONTROLLED;

ALLOCATE FACTOR(NFAC,N);

```

```

      N1 = N;
      L = 1;
      NTAB = 0;
A1:   IF N1 = 0 THEN GO TO A2;
      NTAB = NTAB + N1;
      N1 = N1 - 2;
      GO TO A1;

A2:   ALLOCATE TAB(NTAB),TE(ND2),TE(ND2),TP(ND2);

      DO I = 1 TO N;
        TAB(I) = I;
      END;
      N1 = N;
      N3 = 1;
      DO I = 1 TO ND2;
        TB(I) = N3;
        TP(I) = N3;
        N3 = N3 + N1;
        TE(I) = N3 - 1;
        N1 = N1 - 2;
      END;

      I = 1;
A5:   TP(I) = TP(I) + 1;
      IF TP(I) > TE(I) THEN DO;
        IF I = 1 THEN GO TO A20;
        TP(I) = TE(I);
        I = I - 1;
        GO TO A5;
      END;
      IF I = ND2 THEN GO TO A7;
      K = TB(I+1);
      DO J = TB(I) + 1 TO TE(I);
        IF J = TP(I) THEN GO TO A5;
        TAB(K) = TAB(J);
        K = K + 1;
      END;
A6:   END;
      I = I + 1;
      GO TO A5;

A7:   DO I = 1 TO ND2;
        FACTOR(L,2*I-1) = TAB(TB(I));
        FACTOR(L,2*I) = TAB(TP(I));
      END;
      L = L + 1;
      I = ND2;
      GO TO A5;

A20:  FREE TAB,TB,TE,TP;
      END FORNFAC;

```

```

/* ***** */

```

```

TRANFAC: PROCEDURE;

```

```

/* ROUTINE TO FIND MAPPINGS OF ALL FACTORS BY RELABELING VERTICES */

```

```

DECLARE TEST(N) FIXED BINARY CONTROLLED;
DECLARE SAVE FIXED BINARY;
DECLARE X FIXED (3);

ALLOCATE FTRANS(NFAC,NEDGE),TEST(N);

FTRANS = 0;
IJ = 1;
DO I = 1 TO NFAC;
  DO J = I+1 TO N;
    DO K = 1 TO NFAC;
      IF FTRANS(K,IJ) > 0 THEN GO TO T11;
      TEST = FACTOR(K,*);
      DO L = 1 TO N;
        IF TEST(L) = I THEN M1 = L;
        IF TEST(L) = J THEN DO;
          M2 = L;

```

```

                GO TO T1;
            END;
        END;
T1:            TEST(M1) = J;
                TEST(M2) = I;

                I1 = 1;
T2:            IF TEST(I1) < TEST(I1+1) THEN GO TO T3;
                SAVE = TEST(I1);
                TEST(I1) = TEST(I1+1);
                TEST(I1+1) = SAVE;
T3:            IF I1 = NM1 THEN GO TO T4;
                I1 = I1 + 2;
                GO TO T2;
T4:            I1 = 1;
T5:            DO I2 = I1+2 TO NM1 BY 2;
                IF TEST(I1) < TEST(I2) THEN GO TO T6;
                SAVE = TEST(I1);
                TEST(I1) = TEST(I2);
                TEST(I2) = SAVE;
                SAVE = TEST(I1+1);
                TEST(I1+1) = TEST(I2+1);
                TEST(I2+1) = SAVE;
T6:            END;
                IF I1 = N-3 THEN GO TO T7;
                I1 = I1 + 2;
                GO TO T5;

T7:            I6 = (TEST(2)-2)*NEF;
                I7 = (TEST(2)-1)*NEF;
                I4 = 2;
                I3 = NEF;
T8:            X = I3;
                I3 = CEIL(X/2);
                I6 = I6+I3*(-1)**I4;
                IF I6 > I7 THEN I6 = I7;
                DO I5 = 4 TO N;
                    IF FACTOR(I6,I5) < TEST(I5) THEN DO;
                        I4 = 2;
                        GO TO T10;
                    END;
                    IF FACTOR(I6,I5) > TEST(I5) THEN DO;
                        I4 = 1;
                        GO TO T10;
                    END;
                END;
                FTRANS(I6,IJ) = K;
                FTRANS(K,IJ) = I6;
                GO TO T11;
T10:           IF I3 > 1 THEN GO TO T8;
                PUT PAGE EDIT ('ERROR1') (A);
                GO TO T20;
T11:           END;
                IJ = IJ + 1;
            END;
        END;
T20: END TRANFAC;

```

```

/* *****

```

```
FCRMSCH: PROCEDURE;
```

```

/* THIS PROCEDURE PUTS FACTORS TOGETHER TO FORM ALL SCHEDULES WHICH
   INCLUDE FACTOR 1 */

```

```

DECLARE SCHP(NM1) FIXED BINARY CONTROLLED;
DECLARE POINT(NEDGE) FIXED BINARY CONTROLLED;
DECLARE SVM(NM2) FIXED BINARY CONTROLLED;

```

```

ALLOCATE PCINT(NEDGE);
ALLOCATE SCHP(NM1), SVM(NM2-1);
ALLOCATE SCFEC(MAXSCH,N);

```

```
POINT = 0;
```

```

NCNT = 0;
DO I = 1 TO NM1;
  SCHP(I) = (I-1)*NEF+1;
END;
CO J = 3 TO N BY 2;
  MM = 0;
  DO M5 = 1 TO FACTOR(1,J);
    MM = MM + N - M5;
  END;
  MM = MM - N + FACTOR(1,J+1);
  PCINT(MM) = 1;
END;

I = 2;
F5: IF (SCHP(I)-1)/(I*NEF) < 1 THEN GO TO F7;
IF I = 2 THEN GC TC F20;
F6: SCHP(I) = (I-1)*NEF+1;
I = I - 1;
DO J = 3 TO N BY 2;
  MM = 0;
  DO M5 = 1 TO FACTOR(SCHP(I),J);
    MM = MM + N - M5;
  END;
  MM = MM - N + FACTOR(SCHP(I),J+1);
  POINT(MM) = 0;
END;
SCHP(I) = SCHP(I) + 1;
GO TO F5;

F7: J1 = 0;
DO J = 3 TO N BY 2;
  J1 = J1 + 1;
  MM = 0;
  DO M5 = 1 TO FACTOR(SCHP(I),J);
    MM = MM + N - M5;
  END;
  MM = MM - N + FACTOR(SCHP(I),J+1);
  IF PCINT(MM) = 1 THEN GO TO F8;
  SVMM(J1) = MM;
END;
IF I = NM1 THEN GO TO F9;
DO J1 = 1 TO ND2-1;
  POINT(SVMM(J1)) = 1;
END;
I = I + 1;
GO TO F7;
F8: SCHP(I) = SCHP(I) + 1;
GO TO F5;
F9: IF NCNT = 1000 THEN GO TO F20;
NCNT = NCNT + 1;
DO I1 = 2 TO NM1;
  SCHED(NCNT,I1-1) = SCHP(I1);
END;
F10: I = NM1;
GO TO F6;

F20: END FCRMSCH;

/* ***** */

EQUIV: PROCEDURE;
/* ROUTINE TO FIND EQUIVALENCE CLASSES FOR ALL SCHEDULES WITH
FIXED FACTOR 1. */

DECLARE TRTBL (NEF-1,5) FIXED BINARY CONTROLLED;
ALLOCATE NSCH (NM2);
DO I = 1 TO NCNT;
  SCHED(I,NM1) = I;
  SCHED(I,N) = C;
END;
FLAG = 0; KTRFG = 0; JEXT = 1;

```

```

/* THIS PART EXCHANGES VERTICES OF FACTOR 1 WITHOUT EXCHANGING EDGES. */
JJ = 1;
JJM = NM2;
G1: DO I = 1 TO NCNT;
    DO K = 1 TO NM2;
        NSCH(K) = FTRANS(SCHED(I,K),JJ);
    END;

    CALL FINDSCH;
END;

CALL PRINTEQ;

IF JJM = 0 THEN GO TO G6;
JJ = JJ + 1 + 2*JJM;
JJM = JJM - 2;
GO TO G1;

/* THIS PART EXCHANGES EDGES OF FACTOR 1 */
G6: JJ = 2;
    JJJ = JJ + NM1;
    JJJT = NM1;
    JC = 1;
    JCT = ND2-1;
    KTRFG = 1;

G7: DO I = 1 TO NCNT;
    DO K = 1 TO NM2;
        NSCH(K) = FTRANS(FTRANS(SCHED(I,K),JJ),JJJ);
    END;

    CALL FINDSCH;
END;

CALL PRINTEQ;

IF JCT = 1 THEN GO TO G10;
IF JC = JCT THEN GO TO G9;
JC = JC + 1;
JJ = JJ + 2;
JJJ = JJJ + 2;
GO TO G7;

G9: JCT = JCT - 1;
    JC = 1;
    JJ = JJJ + 2;
    JJJT = JJJT - 2;
    JJJ = JJ + JJJT;
    GO TO G7;

/* THIS PART CREATES A FACTOR 1 FROM OTHER FACTORS. INITIALLY IT
CREATES A TABLE CALLED TRTBL WHICH WILL HOLD ALL TRANSFORMATIONS
NEEDED TO CONVERT A FIRST FACTOR INTO FACTOR 1. */
G10: ALLOCATE TRTBL (2:NEF,5);

    TRTBL = 0;
    KTRFG = 2;
    DO J = 2 TO NEF;
        DO K = NM1 + NM1 TO NEDGE;
            IF FTRANS(J,K) = 1 THEN DO;
                TRTBL(J,1) = K;
                GO TO B1;
            END;
        END;
    END;

B1: END;

JCTR = 1;
B15: JCT = 0;
    DO J = 2 TO NEF;
        IF TRTBL(J,1) = 0 THEN GO TO B3;
        DO JJ = 2 TO NEF;
            IF TRTBL(JJ,JCTR) = 0 THEN GO TO B4;
            DO K = NM1 + NM1 TO NEDGE;

```

```

      IF FTRANS(J,K) = JJ THEN DO:
        TRTEL(J,1) = K:
        J1 = 2:
      B2:   TRTEL(J,J1) = TRTEL(JJ,J1-1):
          IF TRTEL(JJ,J1) = 0 THEN GO TO B3:
          IF J1 = 5 THEN DO:
            PUT EDIT ('ERROR TRTEL') (A):
            GO TO SSEND:
          END:
          J1 = J1 + 1:
          GO TO B2:
        END:
      END:
    B4:   END:
        JCT = JCT + 1:
    B3:   END:

      IF JCT = C THEN GO TO B6:
      JCTR = JCTR + 1:
      IF JCTR = 5 THEN GO TO B15:
      PUT EDIT ('ERROR2 TRTEL') (A):
      GO TO SSEND:

/* THIS PART WILL THEN LINK LATER FACTORS TO A FIRST FACTOR. */
    B6:   DO JFAC = NEF + 1 TO NFAC:
        JF = (JFAC-1)/NEF:
        DO JJ = N TO 2*N-3:
          IF FTRANS(JFAC,JJ) > NEF THEN GO TO B7:
          KK = FTRANS(JFAC,JJ):
          IF KK = 1 THEN GO TO B20:
          GO TO B8:
    B7:   END:
        PUT EDIT ('ERROR3') (A):
        GO TO SSEND:

    B8:   DO I = 1 TO NCNT:
        IF SCHED(1,JF) = JFAC THEN GO TO B12:
        K1 = 1:
        J1 = 1:
        J2 = FTRANS(1,JJ):
    B85:   DO J3 = 1 TO 5:
        IF TRTEL(KK,J3) = 0 THEN GO TO B9:
        J2 = FTRANS(J2,TRTEL(KK,J3)):
        END:
    B9:   NSCH(K1) = J2:
        IF K1 = NM2 THEN GO TO B11:
        IF J1 = JF THEN J1 = J1 + 1:
        J2 = FTRANS(SCHED(1,J1),JJ):
        K1 = K1 + 1:
        J1 = J1 + 1:
        GO TO B85:

    B11:   CALL FINDSCH:
    B12:   END:

      CALL PRINTEG:

    B20: END:

```

```

/* ***** */

```

```

FINDSCH: PROCEDURE:

```

```

/* FIND SCHEDULE IN TABLE AND FORM EQUIVALENCE CLASSES */

```

```

  DECLARE SAVE FIXED BINARY:
  DECLARE X FIXED (3):

  DO K1 = 1 TO NM2-1:
    DO K2 = K1+1 TO NM2:
      IF NSCH(K1) < NSCH(K2) THEN GO TO E4:
      .SAVE = NSCH(K1):
      NSCH(K1) = NSCH(K2):
      NSCH(K2) = SAVE:

```

```

E4:   END:
      END:
      I6 = 0;
      I3 = NCNT;
      I4 = 2;

E5:   X = I3;
      I3 = CEIL(X/2);
      I6 = I6 + I3*(-1)**I4;
      IF I6 > NCNT THEN I6 = NCNT;
      IF I6 < 1 THEN I6 = 1;
      DO I5 = 1 TO NM2;
        IF SCHED(I6,I5) < NSCH(I5) THEN DO:
          I4 = 2;
          GO TO E10;
        END:
        IF SCHED(I6,I5) > NSCH(I5) THEN DO:
          I4 = 1;
          GO TO E10;
        END:
      END:
      IF SCHED(I6,NM1) = SCHED(I,NM1) THEN GO TO E12;
      FLAG = 1;
      IHEAD = SCHED(I,NM1);
      I7 = 1;
E6:   IF SCHED(I7,N) = 0 THEN GO TO E7;
      I7 = SCHED(I7,N);
      GO TO E6;
E7:   SCHED(I7,N) = SCHED(I6,NM1);
      I6 = SCHED(I6,NM1);
E8:   SCHED(I6,NM1) = IHEAD;
      IF SCHED(I6,N) = 0 THEN GO TO E12;
      I6 = SCHED(I6,N);
      GO TO E8;

E10:  IF X > 1 THEN GO TO E5;
      PUT PAGE EDIT ('ERROR SEARCH') (A);
      PUT SKIP(1) EDIT ((NSCH(I5) DO I5 = 1 TO NM2)) (10 F(5));
      GO TO SSEND;

E12:  END FINDSCH;

```

```

/* ***** */

```

```

PRINTEQ: PROCEDURE:

```

```

/* PRINT EQUIVALENCE CLASSES */
IF FLAG = 0 THEN GO TO P20;
PUT PAGE EDIT ('EQUIV. CLASSES AT EXCHANGE',JEXT) (A,F(5));
IF KTRFG = 2 THEN PUT SKIP(1) EDIT ('EXCHANGE OF FACTORS') (A);
IF KTRFG = 0 THEN DO:
  CALL CALCLB;
  PUT SKIP(1) EDIT ('LAEEL CHANGE',L1,'-',L2) (A,F(4),A,F(2));
END:
IF KTRFG = 1 THEN DO:
  CALL CALCLB;
  L3 = L1;
  L4 = L2;
  JJ = JJJ;
  CALL CALCLB;
  PUT SKIP(1) EDIT ('LABEL CHANGES',L3,'-',L4,L1,'-',L2)
    (A,F(4),A,F(2),X(3),F(2),A,F(2));
END:
KCT = 0;
DO I1 = 1 TO NCNT;
  IF SCHED(I1,N) = 0 THEN KCT = KCT + 1;
END:

```

```
PUT SKIP(2) EDIT ('#CLASSES = *.KCT) (A,F(5));
IF KCT = 1 THEN GO TO SSEND:
FLAG = 0:
JEXT = JEXT + 1:

CALCLB: PROCEDURE:
M6 = 1:
DO M5 = 1 TO NM1:
  M6 = M6 + N - M5:
  IF M6 > JJ THEN GO TO P5:
END:
P5: L1 = M5:
  L2 = N + 1 - M6 + JJ:
  END:

P20: END PRINTEQ:

  END EQUIV:

SSEND: END SCHPGM1:
```

```
/* ***** */
```

Computer Program 2

The main program (SCHPGM2) finds representatives from each equivalence class of schedules with 10 teams. It will also work for schedules with 4, 6 and 8 teams.

As a representative from each new equivalence class is generated, a line is printed and a card is punched. The main purpose of the punched output is for restart purposes. The program also prints totals for each equivalence class at the end of every level break for 1-factors from W_2 and W_3 (see Lemmas 3.6, 3.7 and 3.8). Cards are punched at every level break for 1-factors from W_2 . These cards are used as input to the program PUNCHUP.

The program SCHPGM2 reads input cards which contain the following information:

1. N = number of vertices (corresponds to $2n$ in the thesis body).
2. NH = partition on Hamiltonian circuits with the first 1-factor X_1 (corresponds to 2μ in $\Phi(10, \mu)$ of Lemma 3.12).
3. $N\emptyset NH$ = an array of a maximum of 6 possible values of 2μ not to be included in the array $P_{10}(F)$ for any schedule (see Lemma 3.12).
4. $NF2$ = starting value for the 1-factor from W_2 .
If $NF2 = 0$ then the program assumes it is

to generate schedules starting with the first 1-factors in W_2 and W_3 ; i.e., it is not a restart procedure.

5. NF3 = starting value for the 1-factor from W_3 .
6. NNN = estimated number of equivalence classes for this partition $\phi(10, \mu)$.
7. NNIS = number of representative schedules to be read in and stored in PSCHED. NNIS = 0 if it is the first run for the partition.
8. PSCHED = the array where representatives from each equivalence class are stored. If NNIS = 0, then there are no cards to be read in and the program generates all members of PSCHED. If the program is being restarted, any earlier generated representatives are read in by card and stored in PSCHED.

The program generates all 1-factors in the routine TFØRMFC. The routines FACTEST and FTR4TST perform the isomorphism tests described in Lemmas 3.6, 3.7 and 3.8 and store all relevant 1-factors in the array FACSV.

The routine TFØRMSC is the main schedule-producing routine. It calls three subroutines: TCYCLCT, TISØM and PRISØM. The routine TCYCLCT counts k-circuits and forms the arrays CYCLF and CCT which correspond to $P_{10}(F)$ and $Q_4(F)$

(see Chapter 3) respectively. The routine TISØM tests for isomorphism. In particular, TISØM stores characteristic information regarding each equivalence class which can be useful for isomorphism testing. It performs the tests described in Lemma 3.9 parts i and ii. If necessary it calls the subroutine TEQUIV to perform the test from Lemma 3.9 part iii. TEQUIV maps k -circuits onto k -circuits in an order based on low-multiplicity of numbers in $P_{10}(F)$ and $Q_4(F)$. It calls three subroutines to aid it: FILMPC2, MAPCHEK, and ECHECK (which is the final isomorphism check). If a new equivalence class has been found, TISØM calls TMPCIR to find the best combinations of 1-factors to use for mappings of k -circuits for this equivalence class. TMPCIR calls two subroutines, FILMPCC and CHECKS, to help it accomplish this. The purpose of the routine PRISØM is to print and punch the representative of a new equivalence class.

The program PUNCHUP summarizes the cards punched from SCHPGM2. Output from PUNCHUP consists of one card for every equivalence class (396 cards in all for K_{10}) with the following information on it:

1. Card identification: "3".
2. The value of 2μ for the partition $\phi(10, \mu)$.

3. The sequence number as it was generated by SCHPGM2 (within $\Phi(10, \mu)$).
4. The total number of Hamiltonian circuits.
5. The total number of members in the equivalence class with the fixed 1-factor X_1 .
6. Labels for the 1-factors in the representative schedule.

The program SCHLIST reads the cards punched by PUNCHUP and produces the listing given by Appendix A.

The programs were run on the 512K IBM 370-145 at the University of Victoria. They were compiled using the \O S/360 PL/1 F compiler with full optimization ($\text{\O PT} = 2$). Compile time for SCHPGM2 was 80.3 seconds. Storage was allocated dynamically and varied between 102K and 164K according to the value of NNN (the estimated number of equivalence classes for the partition $\Phi(10, \mu)$) read in. Total execution time for all runs to solve the problem of the 1-factorizations of K_{10} added to 770 minutes.

PUNCHUP: PROCEDURE OPTICKS (MAIN):

```

/* THIS PROGRAM TAKES CARD OUTPUT FROM SCHPGM2 AND PUNCHES SUMMARY
CARDS. */

DECLARE PUNCH OUTPUT;
DECLARE CARD1(15) FIXED DECIMAL(5);
DECLARE CARD2(14) FIXED DECIMAL(5);
DECLARE (PCD,TOT) FIXED DECIMAL(5);
OPEN FILE (PUNCH) OUTPUT;
FFG = 0;
CARD2(1) = 3;
A1: GET SKIP(1) EDIT ((CARD1(I) DO I = 1 TO 15)) (15 F(5));
IF FFG = 1 THEN GO TO A2;
FFG = 1;
GO TO A3;
A2: IF CARD1(4) = PCD THEN GO TO A5;
CARD2(3) = CARD2(3) + 1;
CARD2(5) = TGT;
PUT FILE (PUNCH) SKIP(1) EDIT ((CARD2(J) DO J = 1 TO 14)) (4 F(5),X(5),
10 F(5));
IF CARD1(1) = 0 THEN GO TO A10;
IF CARD2(2) = CARD1(2) THEN GO TO A4;
A3: CARD2(3) = 0;
CARD2(2) = CARD1(2);
A4: CARD2(4) = CARD1(5);
DO J = 7 TO 15;
CARD2(J-1) = CARD1(J);
END;
TGT = 0;
PCD = CARD1(4);
A5: IF CARD1(3) = 110 THEN TGT = TGT + CARD1(6)*20;
ELSE TGT = TGT + CARD1(6) * 48;
GO TO A1;
A10: END PUNCHUP;

/* ***** */

```

SCHLIST: PROCEDURE OPTIONS (MAIN):

/* THIS PGM WILL MAKE PROPER LISTINGS OF SCHEDULES FROM CARD INPUT. */

```

DECLARE FACTOR (1000,10) FIXED BINARY;
DECLARE SCHED (400,9) FIXED BINARY;
DECLARE HCTOT(400) FIXED BINARY;
DECLARE AUTCM(400);
DECLARE CYCLF(400,9) FIXED BINARY;
DECLARE CCT(400,10) FIXED BINARY;
DECLARE WGRK(10) FIXED BINARY;
DECLARE CYCT FIXED BINARY;
DECLARE CYC(10) FIXED BINARY;
DECLARE FACWK(1000) FIXED BINARY;
DECLARE DCYCLF(400,9) FIXED BINARY;
DECLARE DCCT(400,10) FIXED BINARY;
DECLARE WK(10) FIXED BINARY;

DECLARE SYSPRINT PRINT;

OPEN FILE (SYSPRINT) PAGESIZE (100);

GET LIST(N);
NM1 = N-1;
  NM2 = N-2;
ND2 = N/2;
NEF = 1;
NC = 1;
SS2: NC = NC+2;
  IF NC > N-2 THEN GO TO SS3;
  NEF = NEF*NC;
  GO TC SS2;
SS3: NFAC = NM1*NEF;
  NTAB = 0;
  N1 = N;
SS4: IF N1 = 0 THEN GO TO SS5;
  NTAB = NTAB + N1;
  N1 = N1-2;
  GO TO SS4;
SS5: PERM = 1;
  DO I = 2 TO ND2;
    PERM = PERM*I;
  END;
  PERM = PERM*NM1*2**ND2;

  NCNT = 0;
  TA = 0;
SS10: NCNT = NCNT+1;
  GET SKIP(1) EDIT (ICD,HCTOT(NCNT),AUTCM(NCNT),(SCHED(NCNT,J) DO J = 1
    TO NM1) (F(5),X(10),F(5),F(5),9 F(5));
  TA = TA + AUTCM(NCNT);
  IF ICD = 3 THEN GO TO SS10;
  NCNT = NCNT-1;

  CALL FORMFAC;

  FACWK = 0;
  DO I = 1 TO NCNT;
    DO J = 1 TC NM1;
      FACWK(SCHED(I,J)) = 1;
    END;
  END;
  I1 = 0;
  DO I = 1 TO NFAC;
    IF FACWK(I) = 1 THEN DO;
      I1 = I1 + 1;
      FACWK(I) = I1;
    END;
  END;

  CALL CYCLECT;

  CALL TORDER;

  CALL PRINT;
  GO TO SS30;

```

```
SS20: PUT PAGE EDIT (*ERROR CARD*,K,CYCT,HCTOT(K)) (A,3 F(5));
```

```
/* ***** */
```

```
FORMFAC: PROCEDURE:
```

```
/* TO GENERATE ALL THE FACTORS FOR N TEAMS */
```

```
DECLARE TAB(NTAB) FIXED BINARY CONTROLLED:
DECLARE (TB(ND2),TE(ND2),TP(ND2)) FIXED BINARY CONTROLLED:
```

```
L = 1:
```

```
A2: ALLOCATE TAB(NTAB),TB(ND2),TE(ND2),TP(ND2):
```

```
DO I = 1 TO N:
  TAB(I) = 1:
```

```
END:
```

```
N1 = N:
```

```
N3 = 1:
```

```
DO I = 1 TO ND2:
```

```
  TB(I) = N3:
```

```
  TP(I) = N3:
```

```
  N3 = N3 + N1:
```

```
  TE(I) = N3 - 1:
```

```
  N1 = N1 - 2:
```

```
END:
```

```
I = 1:
```

```
A5: TP(I) = TP(I) + 1:
```

```
IF TP(I) > TE(I) THEN DO:
```

```
  IF I = 1 THEN GO TO A20:
```

```
  TP(I) = TB(I):
```

```
  I = I - 1:
```

```
  GO TO A5:
```

```
END:
```

```
IF I = ND2 THEN GO TO A7:
```

```
K = TB(I+1):
```

```
DO J = TB(I) + 1 TO TE(I):
```

```
  IF J = TP(I) THEN GO TO A6:
```

```
  TAB(K) = TAB(J):
```

```
  K = K + 1:
```

```
A6: END:
```

```
I = I + 1:
```

```
GO TO A5:
```

```
A7: DO I = 1 TO ND2:
```

```
  FACTOR(L,TAB(TB(I))) = TAB(TP(I)):
```

```
  FACTOR(L,TAB(TP(I))) = TAB(TB(I)):
```

```
END:
```

```
L = L + 1:
```

```
I = ND2:
```

```
GO TO A5:
```

```
A20: FREE TAB,TB,TE,TP:
```

```
END FORMFAC:
```

```
/* ***** */
```

```
CYCLECT: PROCEDURE:
```

```
/* THIS PROCEDURE COUNTS ALL CYCLES BETWEEN ANY TWO FACTORS AND ALSO
COUNTS THE NUMBER OF TIMES EACH FACTOR IS CONTAINED IN EACH TYPE
OF CYCLE. */
```

```
/* ROUTINE COUNTS H.C.'S OR NON-H.C.'S ONLY. */
```

```
/* THIS PROCEDURE ALSO ENUMERATES THE NUMBER OF 4-CYCLES INVOLVING
EACH VERTEX. THIS IS USEFUL FOR CHECKING ISOMORPHISM. */
```

```
CYCLF = 0:
```

```
CCT = 0:
```

```
DO K = 1 TO NCNT:
```

```

CYCT = 0;
DO I = 1 TO NM2;
  DO J = I+1 TO NM1;
    WK = 0;
    I10 = 1;
    JFG = 0;
C1:    I1 = I10;
    ICT = 0;
C2:    I2 = FACTOR(SCHED(K,I),I1);
    ICT = ICT + 1;
    WK(I1) = 1;
    WK(I2) = 1;
    CYC(ICT) = I1;
    I1 = FACTOR(SCHED(K,J),I2);
    ICT = ICT + 1;
    CYC(ICT) = I2;
    IF I1 = I10 THEN GO TO C3;
    GO TO C2;
C3:    IF JFG = 1 THEN GO TO C5;
    IF ICT = N THEN DO:
      CYCT = CYCT + 1;
      CYCLF(K,I) = CYCLF(K,I) + 1;
      CYCLF(K,J) = CYCLF(K,J) + 1;
      GO TO C6;
    END;
    JFG = 1;
C5:    IF ICT = 6 THEN GO TO C7;
    DO I6 = 1 TO 4;
      I7 = CYC(I6);
      CCT(K,I7) = CCT(K,I7) + 1;
    END;
C7:    DO I3 = 1 TO N;
      IF WK(I3) = 0 THEN DO:
        I10 = I3;
        GO TO C1;
      END;
    END;
C6:    END;
C10: END;
  IF CYCT = HCTOT(K) THEN GO TO SS20;
  END;
END CYCLECT;

/* ***** */

TORDER: PROCEDURE:
/* ORDERS H-CIRCUITS ACROSS THE FACTORS AND 4-CYCLES ACROSS THE VERTICES*/
DO I = 1 TO NCNT;
  DO J = 1 TO NM1;
    WK(J) = CYCLF(I,J);
  END;
  J1 = 1;
  MINL = 1;
T5:   DO J = 2 TO NM1;
      IF WK(J) < WK(MINL) THEN MINL = J;
    END;
    OCYCLF(I,J1) = WK(MINL);
    IF J1 = NM1 THEN GO TO T8;
    J1 = J1 + 1;
    WK(MINL) = 99;
    GO TO T5;
T8:   END;

```

```

LINCT = 0;
PUT PAGE EDIT ('REPRESENTATIVE SCHEDULES FROM EACH EQUIVALENCE CLASS FOR *
, 'SCHEDULES OF ORDER', N) (A, A, F(3));

IF IFLG = 0 THEN DO;
  PUT SKIP(2) EDIT ('FACTOR LIST') (A);
  PUT SKIP(2) EDIT ('FCTR', 'EDGES', 'FCTR', 'EDGES')
    (A, X(5), A, X(35), A, X(5), A);
  PUT SKIP(1);
END;

IF IFLG = 1 THEN DO;
  PUT SKIP(2) EDIT ('SCHEDULES BY FACTOR NUMBER') (A);
  PUT SKIP(2) EDIT ('SCHED', '* * * * FACTORS * * * *',
    'ORDER', 'HC', '# HAMILT CIRCUITS ACROSS THE FCTRS',
    '# 4-CIRCUITS ACROSS THE VERTICES')
    (A, X(4), A, X(1), A, X(2), A, X(4), A, X(4), A);
  PUT SKIP(1);
END;

IF IFLG = 2 THEN DO;
  PUT SKIP(2) EDIT ('SCHEDULES WITH CIRCUITS REORDERED FOR EASIER ',
    'REFERENCE - REFER TO PREVIOUS LISTING') (A, A);
  PUT SKIP(2) EDIT ('SCHED', 'HC', '# HAMILT CIRCUITS ACROSS THE FCTRS',
    '# 4-CIRCUITS ACROSS THE VERTICES')
    (A, X(2), A, X(4), A, X(4), A);
  PUT SKIP(1);
END;

END HEAD;

END PRINT;

SS30: END SCHLIST;

/* ***** */

```

SCHPGM2: PROCEDURE OPTICNS (MAIN);

/* THIS PROGRAM FINDS REPRESENTATIVES FROM EACH EQUIVALENCE CLASS OF SCHEDULES WITH 10 TEAMS. IT WILL ALSO WORK FOR SCHEDULES WITH A SMALLER NUMBER OF TEAMS. IT ALSO PERFORMS CERTAIN STEPS TOWARDS THE COUNTING OF THE NUMBER OF MEMBERS IN EACH EQUIVALENCE CLASS. */

DECLARE PUNCH OUTPUT;

DECLARE FACSV (NFAC,NP1) FIXED BINARY CONTROLLED;
 DECLARE NGNH(6);
 DECLARE ISMTB(2,NEF) FIXED BINARY CONTROLLED;
 DECLARE PSCHED(NNN,NM1) FIXED BINARY CONTROLLED;
 DECLARE FF(NFAC) FIXED BINARY CONTROLLED;

OPEN FILE (PUNCH) OUTPUT;

GET LIST (N);
 PUT EDIT ('INPUT INFORMATION') (A);
 PUT SKIP(3) EDIT ('NUMBER OF VERTICES',N) (A,F(5));
 IF N = 10 | N = 8 | N = 6 | N = 4 THEN GO TO SS1;
 GO TO SS10;

SS1: NP1 = N+1;
 NM1 = N-1;
 NM2 = N-2;
 ND2 = N/2;
 N999 = 999;
 NEF = 1;

SS2: NC = NC + 2;
 IF NC > N-2 THEN GO TO SS3;
 NEF = NEF*NC;
 GO TO SS2;

SS3: NFAC = (N-1)*NEF;
 NTAB = 0;
 N1 = N;

SS4: IF N1 = 0 THEN GO TO SS5;
 NTAB = NTAB + N1;
 N1 = N1 - 2;
 GO TO SS4;

SS5: GET LIST (NH);
 PUT SKIP(2) EDIT ('PARTITION ON HAMILTONIAN CIRCUITS WITH FIRST FACTOR',
 ' OF EVERY SCHEDULE',NH) (A,A,F(5));
 IF NH < 0 | NH > NM2 THEN GO TO SS10;
 IF MOD(NH,2) = 1 THEN GO TO SS10;
 GET LIST ((NCNH(J) DO J = 1 TO 6));
 PUT SKIP(2) EDIT ('HAMILTONIAN COUNTS NOT INCLUDED',
 '(NGNH(J) DO J = 1 TO 6)) (A,6 F(5));
 GET LIST (NF2,NF3);
 PUT SKIP(2) EDIT ('STARTING VALUES FOR SCHEDULES') (A);
 PUT SKIP(1) EDIT ('FACTOR 2', NF2) (A,F(5));
 PUT SKIP(1) EDIT ('FACTOR 3', NF3) (A,F(5));

GET LIST (NNN);
 PUT SKIP(2) EDIT ('ESTIMATED NUMBER OF NON-ISOMORPHIC SCHEDULES',NNN)
 (A,F(5));

GET LIST (NNIS);
 PUT SKIP(2) EDIT ('NUMBER OF NON-ISOMORPHIC SCHEDULES READ IN',NNIS)
 (A,F(5));
 ALLOCATE PSCHED(NNN,NM1);
 DO L = 1 TO NNIS;
 GET SKIP(1) EDIT ((PSCHED(L,J) DO J = 1 TO NM1)) (X(15),9 F(5));
 PUT SKIP(1) EDIT ((PSCHED(L,J) DO J = 1 TO NM1)) (9 F(5));
 END;

PUT PAGE EDIT ('FACTOR','EDGES') (A,X(12),A);
 PUT SKIP(1);

CALL TFORMFC;

PUT PAGE EDIT ('LISTING OF SCHEDULES AND FACTOR NUMBERS FOR ORDER',N,
 ' AND PARTITION',NH) (A,F(3),A,F(3));
 PUT SKIP(1);

```

CALL TFORMSC;
PUT PAGE EDIT ('END OF JOB') (A);
GO TO SSEND;
SS10: PUT PAGE EDIT ('INCORRECT DATA') (A);

```

```

/* *****

```

```

TFCRMFC: PROCEDURE:

```

```

/* TO GENERATE ALL THE FACTORS FOR N TEAMS */

```

```

DECLARE TAB(NTAB) FIXED BINARY CONTROLLED;
DECLARE (TB(ND2),TE(ND2),TP(ND2)) FIXED BINARY CONTROLLED;
DECLARE WORK1(N) FIXED BINARY CONTROLLED;
DECLARE WORK3(N) FIXED BINARY CONTROLLED;
DECLARE FCTR(N) FIXED BINARY CONTROLLED;
DECLARE WCRKS(NEF,NP1) FIXED BINARY CONTROLLED;
DECLARE WKSV(2,N) FIXED BINARY CONTROLLED;

```

```

ALLOCATE FACS(NEF,NP1);
ALLOCATE WCRK1(N);
ALLOCATE WORK3(N);
ALLOCATE TAB(NTAB),TB(ND2),TE(ND2),TP(ND2);
ALLOCATE FCTR(N);
ALLOCATE WORKS (NEF,NP1);
ALLOCATE WKSV(2,N);
ALLOCATE ISMTB (2,NEF);
ALLOCATE FF(NFAC);

```

```

LM = 0;
LL = 0;
MFLH = 0;
MFLNH = 0;
MFLG1 = 0;
MFLG2 = 0;
MHC = 0;
MNHC = 0;
IHC = 0;
LM2 = 0;
ISMTB = 0;
DO K2 = 1 TO N BY 2;
  WORK1(K2) = K2 + 1;
  WORK1(K2+1) = K2;
  FACS(1,K2) = K2+1;
  FACS(1,K2+1) = K2;
END;

```

```

DO I = 1 TO N;
  TAB(I) = I;
END;
N1 = N;
N3 = 1;
DO I = 1 TO ND2;
  TB(I) = N3;
  TP(I) = N3;
  N2 = N3 + N1;
  TE(I) = N3 - 1;
  N1 = N1 - 2;
END;

```

```

I = 1;
A5: TP(I) = TP(I) + 1;
IF TP(I) > TE(I) THEN DO;
  IF I = 1 THEN GO TO A20;
  TP(I) = TB(I);
  I = I - 1;
  GO TO A5;
END;
IF I = ND2 THEN GO TO A7;
K = TB(I+1);
DO J = TB(I) + 1 TO TE(I);
  IF J = TP(I) THEN GO TO A6;

```

```

      TAB(K) = TAB(J);
      K = K + 1;
A6:  END;
      I = I + 1;
      GO TO A5;

A7:  DO I = 1 TO ND2;
      FCTR(2*I-1) = TAB(TB(I));
      FCTR(2*I) = TAB(TP(I));
      END;
      LL = LL + 1;
      CALL FACTEST;
      I = ND2;
      GO TO A5;

A20: FREE TAB, TB, TE, TP, FCTR;
      FREE WORK1, WORK3;
      FREE WKSV;
      FREE WORKS;

```

```

/* ***** */

```

FACTEST: PROCEDURE:

```

/* STORES IN AN ARRAY CALLED FACSVM ALL FACTORS WHICH WILL COMBINE
WITH FACTOR 1. DEPENDING ON INPUT DATA ONE OR BOTH OF A
NON-HAMILTONIAN AND A HAMILTONIAN FACTOR 2 ARE PICKED.
FINALLY FACTOR 3'S ARE CHOSEN TO BE NON-ISOMORPHIC WHEN COMBINED
WITH FACTOR 1 AND 2. */

IF FCTR(2) = 2 THEN DO;
  IF LM = 1 THEN GO TO W20;
  LM = 1;
  FF(1) = 1;
  ISOMTB (1,1) = NEF;
  PUT SKIP(1) EDIT (LL, (K3 DO K3 = 1 TO N), 0, NEF)
    (F(4), 5 (X(3), 2 F(4)));
  GO TO W20;
END;

DO K2 = 1 TO N BY 2;
  WORK3(FCTR(K2)) = FCTR(K2+1);
  WORK3(FCTR(K2+1)) = FCTR(K2);
END;
DO K2 = 1 TO N;
  IF WORK3(K2) = WORK1(K2) THEN GO TO W20;
END;

I1 = 2;
ICT = 1;
W3: I1 = WORK3(I1);
    IF I1 = 1 THEN GO TO W5;
    IF MOD(I1, 2) = 0 THEN DO;
      I1 = I1 - 1;
      GO TO W4;
    END;
    I1 = I1 + 1;
W4: ICT = ICT + 1;
    GO TO W3;

W5: IF NH = 0 & ICT = ND2 THEN GO TO W20;
    IF NH = NM2 & ICT = ND2 THEN GO TO W20;
    IF FCTR(2) = 3 THEN DO;
      IF ICT = ND2 THEN DO;
        IF MFLH = 0 THEN DO;
          MFLH = 1;
          LM = LM + 1;
          DO K3 = 1 TO N;
            FACSVM(LM, K3) = WORK3(K3);
            WKSV(LM-1, K3) = FCTR(K3);
          END;
          FACSVM(LM, NP1) = 1;
          FF(LM) = LL;
        END;
        MHC = MHC + 1;
      END;
    END;

```

```

      GO I = 1 TO LM2:
      EQU(2,I) = 0:
      DO J3 = 1 TO N:
      IF WORKS(1,J3)  $\neq$  FACS(1,J3) THEN GO TO D37:
      EQU(J11,I) = 0:
      GO TO D38:
D37:  END:
      EQU(1,I) = 1:
D38:  END:

D30:  WK1 = 0:
      I1 = 1:
      ICTT = 0:

D1:   I3 = 1:
      I1SV = I1:

D2:   I2 = FACS(1,I1):
      FCYC(I3) = I2:
      WK1(I1) = 1:
      WK1(I2) = 1:
      I3 = I3 + 1:
      I1 = FACS(J12,I2):
      FCYC(I3) = I1:
      IF I1 = I1SV THEN GO TO D3:
      I3 = I3 + 1:
      GO TO D2:

D3:   ICTT = ICTT + I3:
      DO I4 = 1 TO N:
      MAP(I4) = I4:
      END:
      IBK = 0:
      J2 = 1:

D5:   DO J1 = 1 TO I3:
      MAP(FCYC(J1)) = FCYC(MOD(J1+J2,I3)+1):
      END:
      GO TO D8:

D14:  J2 = 1:
      IBK = 1:
D15:  J3 = I3 - 2:
      DO J1 = J2 TO J2+I3/2-1:
      J4 = FCYC(J1):
      J5 = FCYC(MOD(J1+J3,I3)+1):
      MAP(J4) = J5:
      MAP(J5) = J4:
      J3 = J3-2:
      END:

D8:   DO I = 1 TO LM2:
      IF EQU(J11,I) = 0 THEN GO TO D7:
      DO J3 = 1 TO N:
      WK2(MAP(J3)) = MAP(WORKS(1,J3)):
      END:
      IF WK2(I)  $\neq$  4 THEN GO TO D7:
      DO J4 = 1 TO LM2:
      DO J5 = 2 TO NM1:
      IF WK2(J5)  $\neq$  WORKS(J4,J5) THEN GO TO D6:
      END:
      IF EQU(J11,I) = EQU(J11,J4) THEN GO TO D7:
      IF EQU(J11,I) < EQU(J11,J4) THEN DO:
      J7 = EQU(J11,I):
      J8 = EQU(J11,J4):
      END:
      ELSE DC:
      J7 = EQU(J11,J4):
      J8 = EQU(J11,I):
      END:
      DO J6 = 1 TO LM2:
      IF EQU(J11,J6) = J8 THEN EQU(J11,J6) = J7:
      END:
      GO TO D7:

D6:   END:
      PUT SKIP(2) EDIT ('ERROR IN FTRATST') (A):
      GO TO SSEND:

```

```

D7:  END;
      IF IBK = 0 THEN DC;
          J2 = J2 + 2;
          IF J2 = 13-1 THEN GO TO D14;
          GC TO D5;
      END;

      IF J2 = 13/2 THEN GO TO D10;
          J2 = J2 + 1;
          GO TO D15;
D10: IF ICTT = N THEN GO TO D34;
      DO I1 = 1 TO N;
          IF WK1(I1) = 0 THEN GC TO D1;
      END;

D34: IF NH = 0 | NF = NP2 THEN GO TO D12;
      IF J11 = 2 THEN GO TO D12;
          J12 = 3;
          J11 = 2;
          DO I = 1 TO LM2;
              DO J3 = 1 TO N;
                  IF WCRKS(I,J3) = FACSV(J12,J3) THEN GO TO D32;
                  EQU(J11,I) = 0;
                  GC TO D33;
              END;
          END;
D32:  END;
          EQU(2,I) = I;
D33:  END;
          GO TO D30;

D12: DO I = 1 TO LM2;
      IF EQU(1,I) = I | EQU(2,I) = I THEN DO;
          KCT1 = 0;
          KCT2 = 0;
          DO J = 1 TO LM2;
              IF EQU(1,J) = I THEN KCT1 = KCT1 + 1;
              IF EQU(2,J) = I THEN KCT2 = KCT2 + 1;
          END;
          LM = LM + 1;
          IF KCT1 > 0 THEN ISGMTB(1,LM) = KCT1;
          IF KCT2 > 0 THEN ISGMTB(2,LM) = KCT2;
          DO K3 = 1 TO NP1;
              FACSV(LM,K3) = WORKS(I,K3);
          END;
          FF(LM) = FF(LM+1);
          K5 = 1;
          DO K4 = 1 TO N;
              IF K4 > WORKS(I,K4) THEN GO TO D13;
              WK3(K5) = K4;
              WK3(K5+1) = WORKS(I,K4);
              K5 = K5+2;
          END;
D13:  END;
          PUT SKIP(1) EDIT (FF(LM),(WK3(K3) DO K3 = 1 TO N), WORKS(I,NP1),
              KCT1,KCT2) (F(4),9 (X(3),2 F(4)));
      END;
      END;

      END FTR4TST;

#20: END FACTEST;
      END TFORMFC;

```

```

/* *****
*/

```

```

TFCRMSC: PROCEDURE;

```

```

/* THIS PROCEDURE PUTS FACTORS TOGETHER TO FORM SCHEDULES. */

```

```

DECLARE SCHP(NM1) FIXED BINARY CONTROLLED;
DECLARE CYCT FIXED BINARY;
DECLARE CYCLF (NM1) FIXED BINARY CONTROLLED;
DECLARE SSCHP(N) FIXED BINARY CONTROLLED;
DECLARE WORK(NM1,N) FIXED BINARY CONTROLLED;
DECLARE (WK(N),CYC(N),CCT(N)) FIXED BINARY CONTROLLED;
DECLARE PWURK (NNN,NM1,N) FIXED BINARY CONTROLLED;

```

```

DECLARE PHAMIL(NNN) FIXED BINARY CONTROLLED;
DECLARE PFACTR(NNN,NM1,2) FIXED BINARY CONTROLLED;
DECLARE FSING(NM1) FIXED BINARY CONTROLLED;
DECLARE FDOUB(NM1) FIXED BINARY CONTROLLED;
DECLARE PVERTX(NNN,N,2) FIXED BINARY CONTROLLED;
DECLARE VSING(N) FIXED BINARY CONTROLLED;
DECLARE VDOUB(N) FIXED BINARY CONTROLLED;
DECLARE CFACTR(0:N,2) FIXED BINARY CONTROLLED;
DECLARE CVERTX(0:NP1,2) FIXED BINARY CONTROLLED;
DECLARE MPCC(N) FIXED BINARY CONTROLLED;
DECLARE MAPP(N) FIXED BINARY CONTROLLED;
DECLARE WTEST(N) FIXED BINARY CONTROLLED;
DECLARE TCYCP(N) FIXED BINARY CONTROLLED;
DECLARE WK1(N) FIXED BINARY CONTROLLED;
DECLARE F3TOT(NNN) FIXED BINARY CONTROLLED;
DECLARE F2TOT(NNN) FIXED BINARY CONTROLLED;
DECLARE PMPCC(NNN,N) FIXED BINARY CONTROLLED;
DECLARE PCL(NNN) FIXED BINARY CONTROLLED;
DECLARE PCFAC(NNN,2) FIXED BINARY CONTROLLED;
DECLARE PSV1(NNN,2) FIXED BINARY CONTROLLED;
DECLARE PSVL1(NNN,2) FIXED BINARY CONTROLLED;
DECLARE PSV2(NNN,2) FIXED BINARY CONTROLLED;
DECLARE PSVL2(NNN,2) FIXED BINARY CONTROLLED;
DECLARE PSYMM(NNN) FIXED BINARY CONTROLLED;

ALLOCATE SCHP(NM1);
ALLOCATE CYCLF (NM1);
ALLOCATE SSCHF(N);
ALLOCATE WORK(NM1,N);
ALLOCATE WK(N),CYC(N),CCT(N);
ALLOCATE PWRK(NNN,NM1,N), PHAMIL(NNN), PFACTR(NNN,NM1,2);
ALLOCATE FSING(NM1),FDOUB(NM1);
ALLOCATE PVERTX(NNN,N,2), VSING(N),VDOUB(N);
ALLOCATE CFACTR(0:N,2), CVERTX(0:NP1,2);
ALLOCATE MPCC(N), MAPP(N), WTEST(N);
ALLOCATE TCYCP(N), WK1(N);
ALLOCATE F3TOT(NNN);
ALLOCATE F2TOT(NNN);
ALLOCATE PMPCC(NNN,N), PCL(NNN), PCFAC(NNN,2);
ALLOCATE PSV1(NNN,2), PSVL1(NNN,2), PSV2(NNN,2), PSVL2(NNN,2);
ALLOCATE PSYMM(NNN);

PHAMIL = N999;
CFACTR(0,1) = -N999;
CFACTR(N,1) = N999;
CVERTX(0,1) = -N999;
CVERTX(NP1,1) = N999;
PSV1 = 0;
PSV2 = 0;
PSVL1 = 0;
PSVL2 = 0;
F3TOT = 0;
F2TOT = 0;
PSYMM = 0;
L = C;
ISOMFG = 0;
NCNT = 0;

IF NNIS = 0 THEN GO TO F2;
INIT = 1;
DO L = 1 TO NNIS;
  DO LJ = 1 TO NM1;
    K2 = PSCHEC(L,LJ);
    SCHP(LJ) = K2;
    DO K1 = 1 TO N;
      WCRK(LJ,K1) = FACS(V(K2,K1));
    END;
  END;
  CALL TCYCLCT;
  CALL TISCH;
  CALL PRISOM;
END;

F2: INIT = 0;
IF NH = 0 | NH = NM2 THEN LFLG = 0;
ELSE IF N = 10 & NH = 4 THEN LFLG = 2;
ELSE LFLG = 1;

```

```

IHC = 0;
IHD = 1;
IF NH = 2 THEN CC:
  IHC = 1;
  IHD = 0;
END;
SCHP(1) = 1;
DO K1 = 1 TO N:
  WCRK(1,K1) = FACSV(1,K1);
END;
I = 2;
DO I1 = 2 TO LM:
  IF FACSV(I1,1) > I THEN DO:
    SCHP(I) = I1;
    SSCHP(I) = I1;
    I = I + 1;
  END;
END;
SSCHP(N) = LM + 1;

IF NF2 = 0 THEN GO TO F30;
SCHP(2) = NF2;
SCHP(3) = NF3;

F30: JSFG = 0;

I = 2;
F5: IF SCHP(I) < SSCHP(I+1) THEN GO TO F7;

IF I = 4 THEN DO:
  FF2 = FF(SCHP(2));
  FF3 = FF(SCHP(3));
  PUT SKIP(2) EDIT ('SUBTOTALS FOR FCTR GRP', FF2,'-',FF3,
    (L1,F3TOT(L1) DO L1 = 1 TO LSV)) (A,F(4),A,F(3),X(6),500(F(3),F(5),
    X(4)));
  MULT = 1SCNTE(SCHP(2)-1,SCHP(3));
  PUT SKIP(1) EDIT ('EXT. TOTS FOR FCTR GRP', FF2,'-',FF3,
    (L1,F3TOT(L1)*MULT DO L1 = 1 TO LSV)) (A,F(4),A,F(3),X(6),
    500(F(3),F(5),X(4)));
  PUT SKIP(2);
  DO L1 = 1 TO LSV:
    F2TOT(L1) = F2TOT(L1) + MULT*F3TOT(L1);
  END;
  F3TOT = 0;
END;

IF I = 3 THEN DO:
  DO L1 = 1 TO LSV:
    LCCT = 0;
    DO L2 = 1 TO NM1:
      IF PFACTOR(L1,L2,1) = NH THEN LCCT = LCCT + 1;
    END;
    F2TOT(L1) = F2TOT(L1) * NM1 / LCCT;
  END;
  FF2 = FF(SCHP(2));
  PUT SKIP(2) EDIT ('COMPLETE TOTALS FOR FACTOR GROUP',FF2) (A,F(4));
  PUT SKIP(1) EDIT ((L1,F2TOT(L1) DO L1 = 1 TO LSV))
    (500 (F(3),F(5),X(4)));
  PUT SKIP(2);
  DO L1 = 1 TO LSV:
    IF F2TOT(L1) = 0 THEN GO TO F55;
    PUT FILE (PUNCH) SKIP(1) EDIT (2,NH,FF2,L1,PHANIL(L1),F2TOT(L1),
    (FF(PSCHEDE(L1,J)) DO J = 1 TO NM1)) (15 F(5));
  END;
  F2TOT = 0;
END;

IF I = 2 THEN GO TO F19;
F6: SCHP(I) = SSCHP(I);
I = I - 1;
K2 = SCHP(I);
IF LFLG = 0 THEN GO TO F37;
IF FACSV(K2,NP1) = IHC THEN JSFG = JSFG - 1;
F37: K2 = K2 + 1;
SCHP(I) = K2;
GO TO F5;

```

```

F7:  K2 = SCHP(I);
     IF I = 3 & ISCNTB(SCHP(2)-1,K2) = 0 THEN GO TO F8;

     IF LFLG = 0 THEN GO TO F38;
     IF LFLG = 1 THEN DO:
       KIHC = FACSV(K2,NP1);
       IF JSFG = 2 & KIHC = IHC THEN GO TO F8;
       IF I < NM2 THEN GO TO F38;
       IF I = NM2 & JSFG = 0 & KIHC = IHD THEN GO TO F9;
       IF I = NM1 & JSFG < 2 & KIHC = IFD THEN GO TO F8;
       GO TO F38;
     END;
     IF I < 6 THEN GO TO F38;
     KIHC = FACSV(K2,NP1);
     IF JSFG = 4 & KIHC = IHC THEN GO TO F8;
     DO J20 = 1 TO 4;
       IF I = N-J20 & JSFG < 5-J20 & KIHC = IHD THEN GO TO F8;
     END;

F38: DO K3 = 2 TO I-1;
     DO K1 = 2 TO NM2;
       IF FACSV(K2,K1) = WORK(K3,K1) THEN GO TO F9;
     END;
     DO K1 = 1 TO N;
       WORK(I,K1) = FACSV(K2,K1);
     END;
     IF LFLG = 0 THEN GO TO F39;
     IF FACSV(K2,NP1) = IHC THEN JSFG = JSFG + 1;
F39: IF I = NM1 THEN GO TO F9;
     I = I + 1;
     GO TO F7;
F8:  SCHP(I) = SCHP(I) + 1;
     GO TO F5;

F9:  CALL TCYCLCT;
     ISOMFG = 0;
     CALL TISCM;
     IF ISOMFG = 1 THEN GO TO F10;
     CALL PRISOM;

F10: I = NM1;
     IF LFLG = 0 THEN GO TO F6;
     IF FACSV(SCHP(NM1),NP1) = IHC THEN JSFG = JSFG - 1;
     GO TO F6;

F19: PUT FILE (PUNCH) SKIP(3);

/* ***** */

TCYCLCT: PROCEDURE;

/* THIS PROCEDURE COUNTS ALL HAMILTONIAN CIRCUITS BETWEEN ANY TWO
1-FACTORS AND ALSO COUNTS THE NUMBER OF TIMES EACH FACTOR IS CONTAINED
IN A HAMILTONIAN CIRCUIT. */
/* THIS PROCEDURE ALSO ENUMERATES THE NUMBER OF 4-CIRCUITS INVOLVING
EACH VERTEX. THIS IS USEFUL FOR CHECKING ISCMORPHISM. */

CYCT = 0;
CYCLF = 0;
CCT = 0;

DO I = 1 TO NM2;
  DO J = I+1 TO NM1;
    WK = 0;
    IIC = 1;
    JFG = 0;
C1:   I1 = I10;
    ICT = 0;
C2:   I2 = WORK(I,I1);
    ICT = ICT + 1;
    WK(I1) = 1;
    WK(I2) = 1;
  
```

```

CYC(ICT) = I1;
I1 = WCRK(J,I2);
ICT = ICT + 1;
CYC(ICT) = I2;
IF I1 = I10 THEN GO TO C3;
GO TO C2;

C3:   IF JFG = 1 THEN GO TO C5;
      IF ICT = N THEN DO:
        CYCT = CYCT + 1;
        CYCLF(I) = CYCLF(I) + 1;
        CYCLF(J) = CYCLF(J) + 1;
        GO TO C6;
      END;
      JFG = 1;

C5:   IF ICT = 6 THEN GO TO C7;
      DO I6 = 1 TO 4;
        I7 = CYC(I6);
        CCT(I7) = CCT(I7) + 1;
      END;

C7:   DO I3 = 1 TO N;
        IF WK(I3) = 0 THEN DO:
          I1C = I3;
          GO TO C1;
        END;
      END;

C6:   END;
      ICYC = CYCLF(I);
      JN = 1;

C8:   IF NONH(JN) > N THEN GO TO C10;
      IF ICYC = NONH(JN) THEN GO TO F10;
      JN = JN + 1;
      GO TO C8;

C10:  END;
      ICYC = CYCLF(NM1);
      JN = 1;

C12:  IF NONH(JN) > N THEN GO TO C20;
      IF ICYC = NONH(JN) THEN GO TO F10;
      JN = JN + 1;
      GO TO C12;

C20:  END TCYCLCT;

/* ***** */

PRISGM:  PROCEDURE;

/*  PRINTING ROUTINE FOR SCHEDULES.  */

      PUT SKIP(2) EDIT ('**',L,'**',(FF{SCHP(I1)} DO I1 = 1 TO NM1))
        (A,F(4),X(3),A, X(7),9 F(5));
      PUT SKIP(1) EDIT {'CIRCT',N,' - ',CYCT,(CYCLF(91) DO I1 = 1 TO NM1))
        (X(1),A,F(4),A,F(3),X(1),9 F(5));
      PUT SKIP(0) EDIT {(CCT(I1) DO I1 = 1 TO N)} (X(58),10 F(5));
      PUT FILE (PUNCH) SKIP(1) EDIT (1,NH,L,(SCHP(I1) DO I1 = 1 TO NM1))
        (12 F(5));

END PRISGM;

/* ***** */

TISOM:  PROCEDURE;

/*  PREPARES SCHEDULE TO BE TESTED FOR ISOMORPHISM.  ALL INFO REGARDING
NON-ISOM SCHEDULES ARE STORED IN TABLES.  */

      I2 = I;
H1:    ISM = 1;
      DO I1 = 2 TO NM1;

```

```

      IF CYCLF(I1) < CYCLF(ISM) THEN ISM = I1;
    END;
    CFACTR(I2,1) = CYCLF(ISM);
    CFACTR(I2,2) = ISM;
    IF I2 = NM1 THEN GO TO H2;
    I2 = I2 + 1;
    CYCLF(ISM) = N999;
    GO TO H1;
H2:  I2 = 1;
H3:  ISM = 1;
    DO I1 = 2 TO N;
      IF CCT(I1) < CCT(ISM) THEN ISM = I1;
    END;
    CVERTX(I2,1) = CCT(ISM);
    CVERTX(I2,2) = ISM;
    IF I2 = N THEN GO TO H4;
    I2 = I2 + 1;
    CCT(ISM) = N999;
    GO TO H3;
H4:  IF INIT = 1 THEN GO TO H12;
    DO L = 1 TO NNN;
      IF PHAMIL(L) = N999 THEN GO TO H10;
      IF CYCT ≠ PHAMIL(L) THEN GO TO H5;
      DO J = 1 TO NM1;
        IF CFACTR(J,1) ≠ PFACTR(L,J,1) THEN GO TO H5;
        IF CVERTX(J,1) ≠ PVERTX(L,J,1) THEN GO TO H5;
      END;
      CALL TEQUIV;
      IF ISCMFG = 0 THEN GO TO H5;
      F3TOT(L) = F3TOT(L) + 1;
      GO TO H20;
H5:  END;
    PUT SKIP(1) EDIT ('*TOO MANY NON-EQUIV SCHEDULES') (A);
    GO TO SSEND;

```

```

/* FORM NEW ISCM. GROUP. */

```

```

H10: DO J = 1 TO NM1;
      PSCHED(L,J) = SCHP(J);
    END;
H12: PHAMIL(L) = CYCT;
    DO J = 1 TO NM1;
      PFACTR(L,J,1) = CFACTR(J,1);
      PFACTR(L,J,2) = CFACTR(J,2);
      PVERTX(L,J,1) = CVERTX(J,1);
      PVERTX(L,J,2) = CVERTX(J,2);
    END;
    PVERTX(L,N,1) = CVERTX(N,1);
    PVERTX(L,N,2) = CVERTX(N,2);

    CALL TMPCIR;

    DO J2 = 1 TO NM1;
      DO J3 = 1 TO N;
        PWCRK(L,J2,J3) = WCRK(J2,J3);
      END;
    END;
    DO I2 = 1 TO NM1;
      CYCLF(CFACTR(I2,2)) = CFACTR(I2,1);
      CCT(CVERTX(I2,2)) = CVERTX(I2,1);
    END;
    CCT(CVERTX(N,2)) = CVERTX(N,1);
    LSV = L;
    IF INIT = 1 THEN GO TO H20;
    F3TOT(L) = 1;

```

```

/* ***** */

```

```

TEQUIV: PROCEDURE;

```

```

/* TESTS FOR EQUIVALENCE BY PREPARING ALL POSSIBLE MAPPINGS. */

```

```

/* BEGINS BY PREPARING CIRCUITS FOR CURRENT SCHED IN MPCC. */

```

```

KK1SV = PCFAC(L,1);

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```

KK2SV = PCFAC(L,2);
IF PSYMM(L) = 0 THEN GO TO R2;

/* FOR SYMMETRIC CASES, SET LIMITS FOR OTHER FACTOR COMBINATIONS
   TO TRY. */

DO K60 = KK2SV+1 TC NMI;
  IF CFACTR(KK2SV,1) = CFACTR(K60,1) THEN GO TO Q30;
  K65 = K60 - 1;
  GO TO Q32;
Q30: END;
K65 = NMI;
Q32: IF KK1SV = 1 & CFACTR(1,1) = CFACTR(2,1) THEN DO;
  K72 = KK1SV;
  K73 = K65-1;
  END;
  ELSE DO;
  K72 = KK1SV;
  K73 = KK1SV;
  END;
  K74 = KK2SV;
  K75 = K65;

R2: K1 = CFACTR(KK1SV,2);
K2 = CFACTR(KK2SV,2);

R3: KFG = 0;
CALL FILMPC2;
IF KFG = 1 & PCL(L) = N THEN GO TO R5;
IF KFG = 0 & PCL(L) = 4 THEN GO TO R20;
GO TC R80;

/* CIRCUIT IS HAMILTONIAN. */

R5: K8 = N;
K7 = 1;
IF PSV1(L,1) = 0 THEN DO;
  K20 = 1;
  DO K22 = 1 TC N;
    DO KFLG = 0 TO 1;
      CALL MAPCHEK;
      CALL ECHECK;
    END;
  END;
  GO TO R80;
  END;
  K22 = PSV1(L,1);
  K4 = CVERTX(PSV1(L,1),2);
  DO K5 = 1 TC N;
    IF MPCC(K5) = K4 THEN GO TO R6;
    K20 = K5;
    GO TO R7;
  END;
R6: END;
R7: DO KFLG = 0 TO 1;
  CALL MAPCHEK;
  CALL ECHECK;
  END;
  IF PSV1(L,2) = 0 THEN GO TO R80;
  IF K22 = PSV1(L,2) THEN GO TO R80;
  K22 = PSV1(L,2);
  GO TO R7;

/* CIRCUITS ARE NOT HAMILTONIAN. */

R20: IF PSV1(L,1) > 0 THEN DO;
  K4S = PSV1(L,1);
  K4F = K4S;
  K4 = CVERTX(PSV1(L,1),2);
R21: DO K5 = 1 TO 4;
  IF MPCC(K5) = K4 THEN GO TO R22;
  K4F1 = K5;
  GO TC R25;
R22: END;
  IF K = 8 THEN DO;
  DO K5 = 1 TO 4;
    SAVE = MPCC(K5);
    MPCC(K5) = MPCC(K5+4);
  
```

```

        MPCC(K5+4) = SAVE;
        END;
        GC TO R21;
    END;
    IF PSV1(L,2) = 0 THEN GO TO R80;
    IF K4 = CVERTX(PSV1(L,2),2) THEN GO TO R80;
    K4 = CVERTX(PSV1(L,2),2);
    GC TO R21;
END;

ELSE DO:
    K4S = 1;
    K4F = 4;
    KF1 = 1;
END;

R25: IF PSV2(L,1) > 0 THEN DO:
    K6S = PSVL2(L,1);
    K6F = K6S;
    K4 = CVERTX(PSV2(L,1),2);
R27: DO K5 = 5 TO N;
    IF MPCC(K5) = K4 THEN GO TO R26;
    KF2 = K5-4;
    GC TO R28;
R26: END;
    IF PSV2(L,2) = 0 THEN GO TO R80;
    IF K4 = CVERTX(PSV2(L,2),2) THEN GO TO R80;
    K4 = CVERTX(PSV2(L,2),2);
    GO TO R27;
END;

ELSE DO:
    K6S = 1;
    K6F = 6;
    IF N = 8 THEN K6F = 4;
    KF2 = 1;
END;

R28: KL8 = 6;
    IF N = 8 THEN KL8 = 4;

/* ACTUAL MAPPING AND EQUIV TESTING FOR NON-HAMILT CIRCUIT. */
R30: DO K34 = K4S TO K4F;
    DO KFLG1 = 0 TO 1;
        K8 = 4;
        KFLG = KFLG1;
        K20 = KF1;
        K22 = K34;
        K7 = 1;
        CALL MAPCHEK;
        K8 = KL8;
        K20 = KF2;
        K7 = 5;
        DO K22 = K6S TO K6F;
            DO KFLG = 0 TO 1;
                CALL MAPCHEK;
                CALL ECHECK;
            END;
        END;
    END;
END;

IF K6S = PSVL2(L,2) THEN GO TO R35;
IF PSV2(L,2) = 0 THEN GO TO R35;
K6S = PSVL2(L,2);
K6F = K6S;
GO TO R30;

R35: IF K4S = PSVL1(L,2) THEN GO TO R80;
    IF PSV1(L,2) = 0 THEN GO TO R80;
    K4S = PSVL1(L,2);
    K4F = K4S;
    IF PSV2(L,1) = 0 THEN GO TO R30;
    K6S = PSVL2(L,1);
    K6F = K6S;
    GO TO R30;

```

```

/* SCHEDULES ARE NOT ISOMORPHIC. */
R80: ISCMFG = 0;
IF PSYMM(L) = 0 THEN GO TO R100;

/* FOR SYMMETRIC SCHEDULES, TRY ANOTHER FACTOR CO4BINATION. */
IF K74 = K75 THEN DO:
  IF K72 = K73 THEN GO TO R100;
  K72 = K72 + 1;
  K74 = K72;
END;
K74 = K74 + 1;
K1 = CFACTR(K72,2);
K2 = CFACTR(K74,2);
GO TO R3;

/* SCHEDULES ARE ISCMORPHIC. */
R90: ISCMFG = 1;

/* ***** */

FILMPC2: PROCEDURE;

/* FILLS THE ARRAY MPCC WITH THE CIRCUITS THAT MATCH THE CIRCUITS SAVED
IN PMPCC. */

WK1 = 0;
I1P = 1;
K10 = 0;
T1: I1 = I1P;
ICTP = 0;
T2: I2 = WORK(K1,I1);
ICTP = ICTP + 2;
TCYCP(ICTP-1) = I1;
TCYCP(ICTP) = I2;
WK1(I1) = 1;
WK1(I2) = 1;
I1 = WORK(K2,I2);
IF I1 = I1P THEN GO TO T5;
GO TO T2;
T5: IF ICTP = N THEN GO TO T6;
K10 = K10 + 1;
IF N = 8 THEN DO:
  IF K10 = 1 THEN K6 = 1;
  ELSE K6 = 5;
END;
IF N = 10 THEN DO:
  IF ICTP = 4 THEN K6 = 1;
  ELSE K6 = 5;
END;
GO TO T7;
T6: K6 = 1;
KFG = 1;
T7: DO K5 = K6 TO K6+ICTP-1;
MPCC(K5) = TCYCP(K5-K6+1);
END;
IF KFG = 1 | K10 = 2 THEN GO TO T15;
DO K5 = 1 TO N;
IF WK1(K5) = 1 THEN GO TO T8;
I1P = K5;
GO TO T1;
T8: END;
T15: END FILMPC2;

/* ***** */

MAPCHK: PROCEDURE;

/* PREPARES THE MAPPINGS FROM THE CIRCUITS. */
DO K5 = 0 TO K8-1;

```

```

K9 = KFLG * K5 * 2;
MAPP(MPCC(MOD(K5+K20-1,K8)+K7)) = MPCC(L,MOD(K5+K22-1-K9,K8)+K7);
END;
END MAPCHEK;

```

```

/* ***** */

```

ECHECK: PROCEDURE;

```

/* THE FINAL ISOMORPHISM CHECK. IF ISOMORPHISM OCCURS, ROUTINE
BRANCHES TO STATEMENT R90 OF TEQUIV. */

DO J = 1 TO NM2;
  DO K = 1 TO N;
    WTEST(MAPP(K)) = MAPP(WORK(J,K));
  END;
  KK = WTEST(1) - 1;
  DO K = 2 TO NM1;
    IF PWORX(L, KK, K) = WTEST(K) THEN GO TO Z20;
  END;
END;
GO TO R90;
Z20: END ECHECK;

R100: END TEQUIV;

```

```

/* ***** */

```

TMPDIR: PROCEDURE;

```

/* THIS PROCEDURE FINDS THE BEST CIRCUIT FOR FUTURE ISOM MAPPINGS AND
SAVES ASSOCIATED INFORMATION. */

FSING = 0;
FOOUB = 0;
VSING = 0;
VDOUB = 0;

/* THIS PART FINDS SINGLETON AND DOUBLETION FACTORS AND VERTICES. */

J1 = 0;
DO J3 = 1 TO NM1;
  IF CFACTR(J3-1,1) = CFACTR(J3,1) & CFACTR(J3,1) = CFACTR(J3+1,1)
  THEN DO;
    J1 = J1 + 1;
    FSING(J1) = J3;
  END;
END;
J2 = 0;
DO J3 = 1 TO NM2;
  IF CFACTR(J3-1,1) = CFACTR(J3,1) & CFACTR(J3,1) = CFACTR(J3+1,1)
  & CFACTR(J3+1,1) = CFACTR(J3+2,1) THEN DO;
    J2 = J2 + 2;
    FDOUB(J2-1) = J3;
    FDOUB(J2) = J3+1;
  END;
END;
J1 = 0;
DO J3 = 1 TO N;
  IF CVERTX(J3-1,1) = CVERTX(J3,1) & CVERTX(J3,1) = CVERTX(J3+1,1)
  THEN DO;
    J1 = J1 + 1;
    VSING(J1) = J3;
  END;
END;
J2 = 0;
DO J3 = 1 TO NM1;
  IF CVERTX(J3-1,1) = CVERTX(J3,1) & CVERTX(J3,1) = CVERTX(J3+1,1)
  & CVERTX(J3+1,1) = CVERTX(J3+2,1) THEN DO;
    J2 = J2 + 2;
    VDOUB(J2-1) = J3;
    VDOUB(J2) = J3+1;
  END;
END;

```

```

      END:
    END:

    KFG = 0:
    MPCC = 0:

/* FINDS CIRCUITS BETWEEN SINGLETON FACTORS. */
    DO JS = 1 TO NM1:
      IF FSING(JS) = 0 THEN DO:
        J4 = JS - 1:
        GO TO Y1:
      END:
    END:
    J4 = NM1:

Y1: DO J1 = 1 TO J4 - 1:
      KK1 = FSING(J1):
      K1 = CFACTR(KK1,2):
      DO J2 = J1+1 TO J4:
        KK2 = FSING(J2):
        K2 = CFACTR(KK2,2):
        CALL FILMPCC:
        IF KFG = 1 THEN GO TO Y5:
        IF KFG = 0 THEN CALL CHECKS:
      END:
    END:

/* FIND CIRCUITS BETWEEN DOUBLETEN FACTORS. */
    DO JS = 1 TO NM1 BY 2:
      IF FDOUB(JS) = 0 THEN DO:
        J4 = JS - 2:
        GO TO Y2:
      END:
    END:

Y2: DO J1 = 1 TO J4 BY 2:
      KK1 = FDCUB(J1):
      KK2 = FDCUB(J1+1):
      K1 = CFACTR(KK1,2):
      K2 = CFACTR(KK2,2):
      CALL FILMPCC:
      IF KFG = 1 THEN GO TO Y5:
      IF KFG = 0 THEN CALL CHECKS:
    END:

/* FOR SYMMETRIC CASES NOT COVERED BEFORE */
    IF MPCC(1) = 0 THEN DO:
      IF FSING(1) = 0 THEN DO:
        KK1 = 1:
        KK2 = 2:
      END:
      ELSE DO:
        KK1 = FSING(1):
        KK2 = 1:
        IF KK1 = KK2 THEN KK2 = 2:
      END:
      K1 = CFACTR(KK1,2):
      K2 = CFACTR(KK2,2):
      PSYMM(L) = 1:
      CALL FILMPCC:
      IF KFG = 0 THEN CALL CHECKS:
    END:

/* FILL IN ARRAYS */
Y5: DO K5 = 1 TO N:
      PMPCC(L,K5) = MPCC(K5):
    END:
    IF KFG = 1 THEN PCL(L) = N:
    ELSE PCL(L) = 4:
    PCFAC(L,1) = KK15V:
    PCFAC(L,2) = KK25V:
    IF KFG = 1 THEN DO:
      IF VSING(1) = 0 THEN GO TO Y7:

```

```

PSV1(L,1) = VSING(1);
K9 = CVERTX(VSING(1),2);
DO K5 = 1 TO N;
  IF MPCC(K5) = K9 THEN GO TO Y6;
  PSVL1(L,1) = K5;
  GO TO Y20;
END;
Y5: IF VDCUB(1) = 0 THEN GO TO Y20;
Y7: PSV1(L,1) = VDCUB(1);
PSV1(L,2) = VDCUB(2);
K8 = CVERTX(VDCUB(1),2);
K9 = CVERTX(VDCUB(2),2);
DO K5 = 1 TO N;
  IF MPCC(K5) = K8 THEN PSVL1(L,1) = K5;
  IF MPCC(K5) = K9 THEN PSVL1(L,2) = K5;
END;
GO TO Y20;
END;
IF K51 > 0 THEN DO;
  PSV1(L,1) = K51S;
  PSV1(L,2) = K53S;
  PSVL1(L,1) = K51;
END;
ELSE IF K53 > 0 THEN DO;
  PSV1(L,1) = K53S;
  PSV1(L,2) = K54S;
  PSVL1(L,1) = K53;
  PSVL1(L,2) = K54;
END;
IF K52 > 0 THEN DO;
  PSV2(L,1) = K52S;
  PSV2(L,2) = K55S;
  PSVL2(L,1) = K52 - 4;
END;
ELSE IF K55 > 0 THEN DO;
  PSV2(L,1) = K55S;
  PSV2(L,2) = K56S;
  PSVL2(L,1) = K55 - 4;
  PSVL2(L,2) = K56 - 4;
END;
END;

/* ***** */

FILMPCC: PROCEDURE;
/* FILLS THE ARRAY MPCC WITH THE BEST CIRCUITS BETWEEN 2 FACTORS. */

WK1 = 0;
I1P = 1;
K10 = 0;
T1: I1 = I1P;
ICTP = 0;
T2: I2 = WCRK(K1, I1);
ICTP = ICTP + 2;
TCYCP(ICTP-1) = I1;
TCYCP(ICTP) = I2;
WK1(I1) = 1;
WK1(I2) = 1;
I1 = WCRK(K2, I2);
IF I1 = I1P THEN GO TO T5;
GO TO T2;
T5: IF ICTP = N THEN GO TO T6;
/* KFG = 2 MEANS THAT THIS CIRCUIT CANNOT BE BETTER THAN AN EARLIER ONE */
IF KFG = 2 THEN GO TO T5;
K10 = K10 + 1;
IF N = 8 THEN DO;
  IF K10 = 1 THEN K6 = 1;
  ELSE K6 = 5;
END;
IF N = 10 THEN DO;
  IF ICTP = 4 THEN K6 = 1;
  ELSE K6 = 5;
END;
GO TO T7;
T6: K6 = 1;

```

```

KFG = 1;
T7: DO K5 = K6 TC K6+ICTP-1;
      MPCC(K5) = TCYCF(K5-K6+1);
      END;
      IF KFG = 1 | K10 = 2 THEN GO TO T10;
      DO K5 = 1 TC N;
        IF WK1(K5) = 1 THEN GO TO T8;
        IIP = K5;
        GO TO T1;
T8: END;
/* SAVES FACTORS IN CASE THIS CIRCUIT TURNS OUT TO BE THE BEST */
T10: KK1SV = KK1;
      KK2SV = KK2;
T15: END FILMPCC;

```

```

/* ***** */

```

CHECKS: PROCEDURE:

```

/* THIS ROUTINE CHECKS NON-HAMIL. CIRCUITS FOR SINGLETON AND DOUBLETON
VERTICES AND STORES THEIR LOCATIONS. */
K51 = 0;
K52 = 0;
K53 = 0;
K54 = 0;
K55 = 0;
K56 = 0;
K53S = 0;
K55S = 0;
DO K6 = 1 TC N;
  K4 = VSING(K6);
  IF K4 = 0 THEN GO TO U6;
  K9 = CVERTX(K4,2);
  IF K51 > 0 THEN GO TO U2;
  DO K5 = 1 TC 4;
    IF MPCC(K5) = K9 THEN GO TO U1;
    K51S = K4;
    K51 = K5;
    GO TO U4;
U1: END;
U2: IF K52 > 0 THEN GO TO U4;
    DO K5 = 5 TO N;
      IF MPCC(K5) = K9 THEN GO TO U3;
      K52S = K4;
      K52 = K5;
      GO TO U4;
U3: END;
U4: IF K52*K51 = 0 THEN GO TO U10;
    END;
U6: DO K6 = 1 TC N BY 2;
      K3 = VDCUB(K6);
      IF K3 = 0 THEN GO TO U20;
      K4 = VDQLH(K6+1);
      K8 = CVERTX(K3,2);
      K9 = CVERTX(K4,2);
      IF K51 > 0 THEN GO TO U7;
      K53 = 0;
      K54 = 0;
      K53S = 0;
      DO K5 = 1 TC 4;
        IF MPCC(K5) = K8 THEN DO;
          K53 = K5;
          K53S = K3;
        END;
        IF MPCC(K5) = K9 THEN DO;
          K54 = K5;
          K54S = K4;
        END;
      END;
      IF K53 > 0 & K54 = 0 THEN DO;
        K51 = K53;
        K51S = K53S;

```

```

      KS3S = K4;
END;
IF KS4 > 0 & K53 = 0 THEN DO:
  KS1 = KS4;
  KS1S = KS4S;
  KS3S = K3;
END;
U7:  IF KS2 > 0 THEN GO TO U8:
      K55 = 0;
      K56 = 0;
      K55S = C;
      DC K5 = 5 TO N;
      IF MPCC(K5) = K8 THEN DC:
        K55 = K5;
        K55S = K3;
      END;
      IF MPCC(K5) = K9 THEN DO:
        K56 = K5;
        K56S = K4;
      END;
      END;
      IF K55 > 0 & K56 = 0 THEN DC:
        K52 = K55;
        K52S = K55S;
        K55S = K4;
      END;
      IF K56 > 0 & K55 = 0 THEN DO:
        K52 = K56;
        K52S = K56S;
        K55S = K3;
      END;
U8:  IF K52*K51 = 0 THEN GO TO U10:
      END;
      GO TO U20;
U10: KFG = 2;
U20: END CHECKS;

Y20: END TPCIR;

H20: END TISGM;

F20:  END TFORMSC;

SEND:  END SCHPGM2;

```

```

/* ***** */

```

VITA

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ON 1-FACTORIZATIONS OF THE COMPLETE
GRAPH AND THE RELATIONSHIP TO
ROUND ROBIN SCHEDULES

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July 1973
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